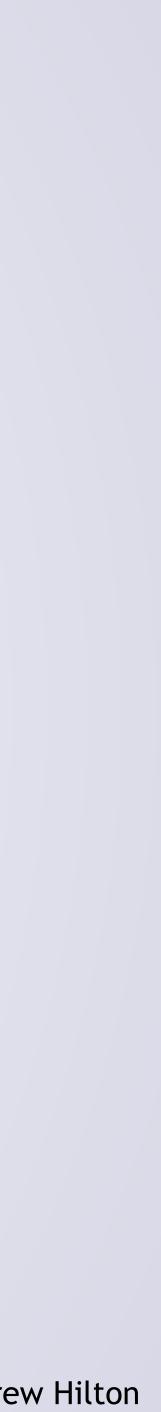
Engineering Robust Server Software Exceptions



Brian Rogers Duke ECE Used with permission from Drew Hilton





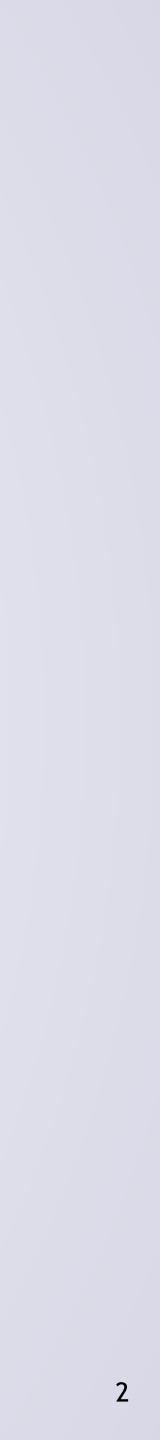
- Handling problems: exceptions
- C++
 - temp-and-swap
 - RAII
 - Smart Pointers
- Java
 - finally •
 - specifications
 - finalizers (and why they are not what you need for this)



Exceptions

C++

Java

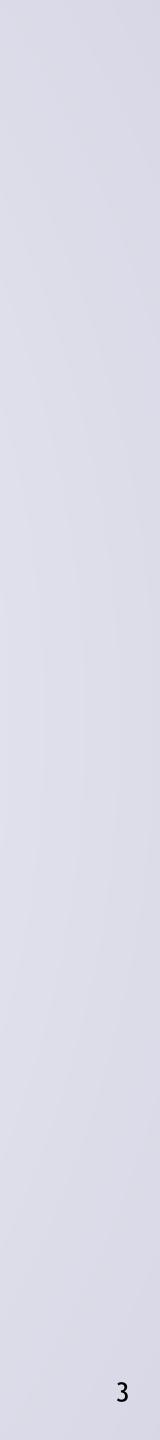




- Review: exceptions = way to handle problems
 - Thing goes wrong? throw exception •
 - Know how to deal with problem? try/catch exception •
 - In python, try/except
- Why exceptions?
 - Return error code? Cluttery, easy to forget/ignore
 - Do nothing? Automatically pass problem to caller •
 - Provide details about error



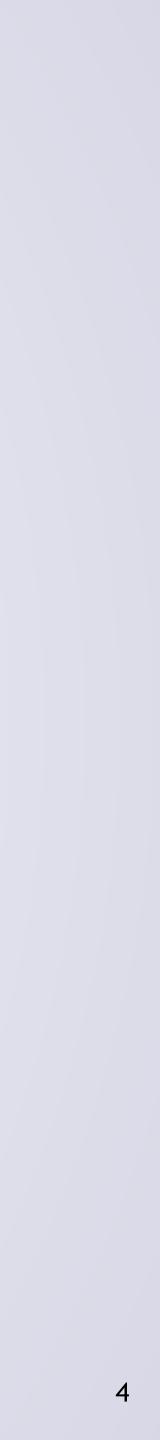
Exceptions



Exceptions: Downsides

- So exceptions: best idea ever?
- Downsides too
 - Unexpected things happen in code
 - Well, that is true anyways
 - Used improperly: corrupted objects, resource leaks, ...
- Bottom line:
 - Good if you do all things right







- Continued review: exception safety
 - Remind us of the four levels of exceptions safety? •

	No Throw	Will not throw a any exceptions
uarantees 	Strong	No side-effects unmodified, and
5	Basic	Objects remain invariants rema
	None	Does not provid Unacceptable in
Stronger Gua	Basic	Objects rema invariants rem Does not prov



Exception Safety

- any exception. Catches and handles throw by operations it uses if an exception is thrown: objects are d no memory is leaked in valid states: no dangling pointers, ain intact. No memory is leaked
- de even a basic exception guarantee. n professional code.







LList & operator=(const LList & rhs) { if (this != &rhs) { deleteAll(); Node * curr = rhs.head; while (curr != null) { addToBack(curr->data); curr = curr->next; return *this;



Exception Safety

- Which guarantee does this make? A: Strong
- **B:** Basic
- **C:** No Guarantee
- D: Need more info...







LList & operator=(const LList & rhs) { if (this != &rhs) { deleteAll(); Node * curr = rhs.head; while (curr != null) { addToBack(curr->data); curr = curr->next; return *this;



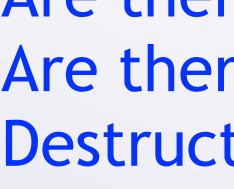
Exception Safety

Which guarantee does this make? - Need to know what guarantees these make!





void deleteAll() { while (head != nullptr) { Node * temp = head->next; B: Strong delete head; head = temp; tail = nullptr;





Exception Safety

Which guarantee does deleteAll() make? A: No Throw C: Basic **D:** No Guarantee

Are there any function calls here? Are there any hidden calls? Yes, the destructor Destructors should always be no-throw



Exception Safety

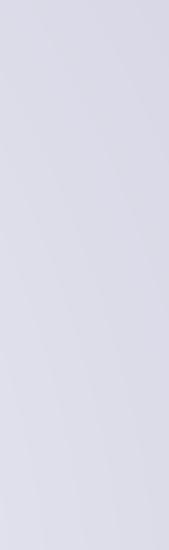
template<typename T> class LList { //other things omitted, but typical

void addToBack(const T& d) { if (tail == nullptr) { head = tail = newNcelse { tail->next = newNod newNode->prev = tai tail = newNode;



- Which guarantee does addToBack() make?
- Depends on copy constructor for T Node * newNode = new Node(d, nullptr, tail) < Could throw memory but does so bef

ode;	T's Copy Constructor	addToBack()
مام	No Throw	Strong
de; il;	Strong	Strong
	Basic	Basic
	No Guarantee	No Guarantee







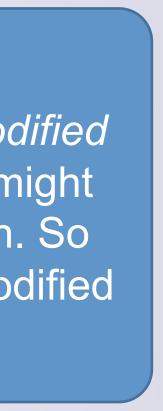
Why? LList & operator=(const LList & rhs) { The list is *being modified* //no throw if (this != &rhs) { when addToBack might //no throw throw an exception. So deleteAll(); we'd leave list in modified Node * curr = rhs.head; //no throw state. while (curr != null) { //no throw addToBack(curr->data); //strong [let us suppose] //no throw curr = curr->next; //no throw return *this;



Exception Safety

Which guarantee does this make? Basic











C++

____st & operator=(const LList & rhs) { if (this != &rhs) { Node * temp = rhs.head; Node * n1 = nullptr; Node * n2 = nullptr; if (temp != nullptr) { temp = temp->next; while (temp != null) { n2 = n2 - next;temp = temp->next; } deleteAll(); head = n1; tail = n2; No guarantee! :(return *this;

An attempt at improving safety: making new temp list before deleting the old one.

n1 = n2 = new Node(temp->data, nullptr, nullptr); n2->next = new Node(temp->data, n2, nullptr);

Which guarantee does this version make?

Why? If we have exception *while building* the new list, then that memory is lost and therefore leaked.



LList & operator=(const LList & rhs) { if (this != &rhs) { LList temp(rhs) < Stack allocated! std::swap(temp.head, head); std::swap(temp.tail, tail);

temp is auto-deleted here

return *this;



Exception Safety

A good strategy to improve exception safety: temp-andswap (also called copy-andswap)

Which guarantee does this make?

Strong!

Why?

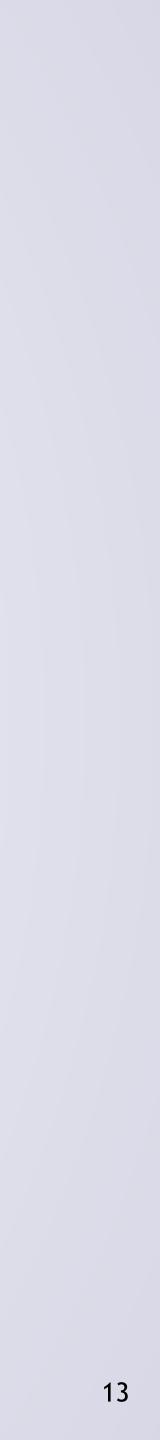
Only place we can get an exception is when making temp (no changes made yet); swap is no-throw.



Temp-and-swap

- Common idiom for strong guarantees: temp-and-swap
 - Make temp object
 - Modify temp object to be what you want this to be
 - swap fields of temp and this
 - temp destroyed when you return (destructor cleans up state) •
 - Exception? temp destroyed in stack unwinding
- Downside?
 - Change only some state: may be expensive to copy entire object





What About This Code...

template<typename T> class LList { //other things omitted, but typical

LList & operator=(const LList & rhs) {

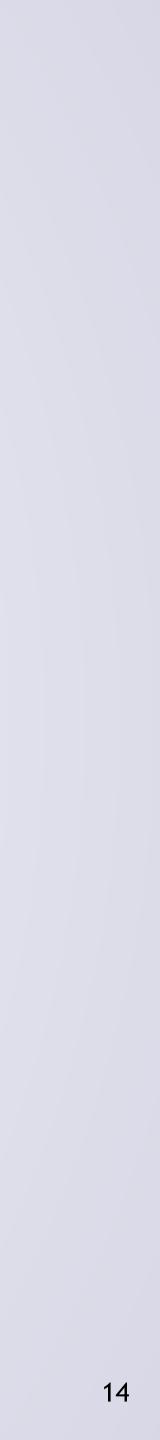
if (this != &rhs) { m.lock(); rhs.m.lock(); std::swap(temp.head, head); std::swap(temp.tail, tail); rhs.m.unlock(); m.unlock();

return *this;



LList temp(rhs); //What if this throws?

We've acquired locks and are not releasing them!





LList & operator=(const LList & rhs) {

if (this != &rhs) { LList temp(rhs); std::swap(temp.head, head); std::swap(temp.tail, tail);

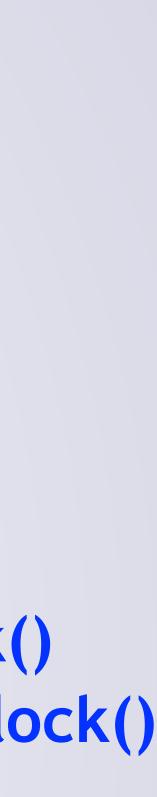
return *this;



Duke

How About Now?

- std::lock guard<std::mutex> lck1(m); //calls m.lock()
- std::lock guard<std::mutex> lck2(rhs.m);//calls rhs.m.lock()





LList & operator=(const LList & rhs) {

if (this != &rhs) { std::lock guard<std::mutex> lck1(m); LList temp(rhs); A: std::swap(temp.head, head); std::swap(temp.tail, tail); **B:** }

C: return *this;

D: They are not unlocked anywhere

};

How About Now?

Where are these locks unlocked?

- std::lock guard<std::mutex> lck2(rhs.m);



LList & operator=(const LList & rhs) {

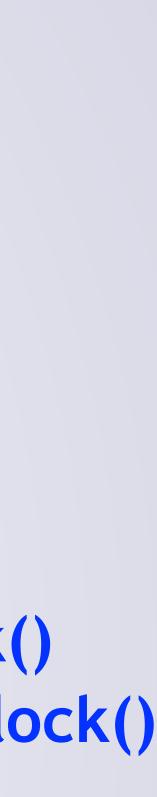
if (this != &rhs) { LList temp(rhs); std::swap(temp.head, head); std::swap(temp.tail, tail); } //destructor of lock_guard calls .unlock() return *this;



Duke

How About Now?

- std::lock guard<std::mutex> lck1(m); //calls m.lock()
- std::lock guard<std::mutex> lck2(rhs.m);//calls rhs.m.lock()





LList & operator=(const LList & rhs) {

if (this != &rhs) { std::lock guard<std::mutex> lck1(m); std::lock guard<std::mutex> lck2(rhs.m); LList temp(rhs); std::swap(temp.head, head); std::swap(temp.tail, tail); Locks are auto-released here (either naturally or on exception)

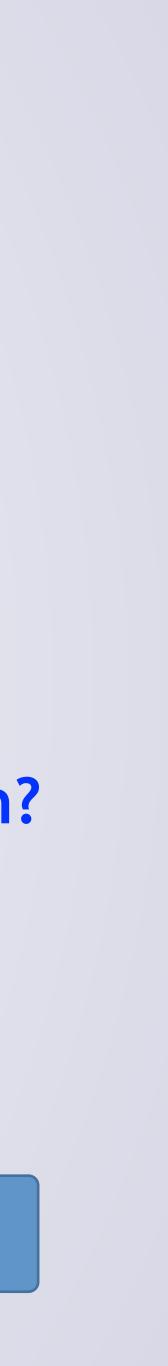
return *this;



How About Now?

//what if exn?

This is an example of **RAII**...



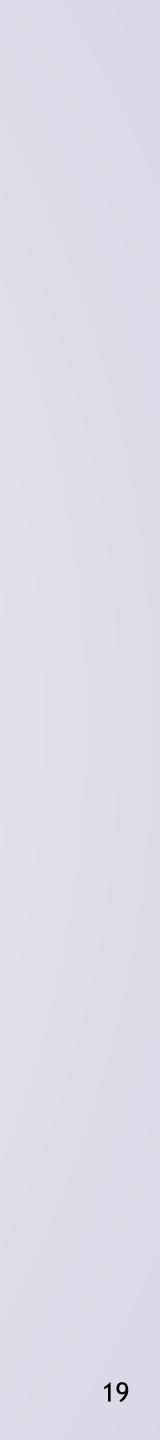
Resource Acquisition Is Initialization •

- Resource lifetime tied to object lifetime •
- Allocation during initialization •
- Released during destruction
- Example resources:
 - Mutex: lock/unlock
 - Heap Memory: new/delete •
 - File: open/close •
- Exception safety benefits?

Release-on-destruction means we release resources whether we go down the normal execution path or exception path.



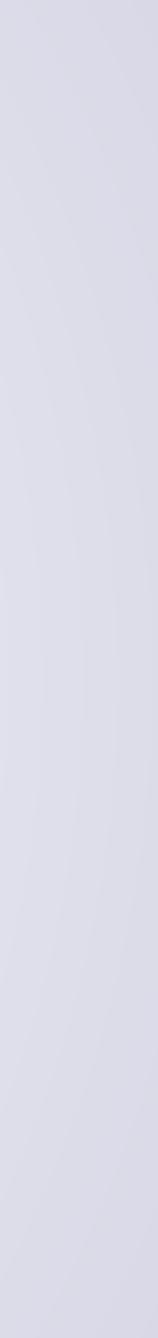




RAll with Heap Objects

- "Smart Pointers"
 - Objects that wrap pointer and provide RAII
 - C++03: std::auto_ptr (deprecated)
- C++11:
 - std::unique_ptr
 - std::shared_ptr
 - std::weak_ptr





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std::unique ptr<Thing> thing1 (new Thing); //other code here

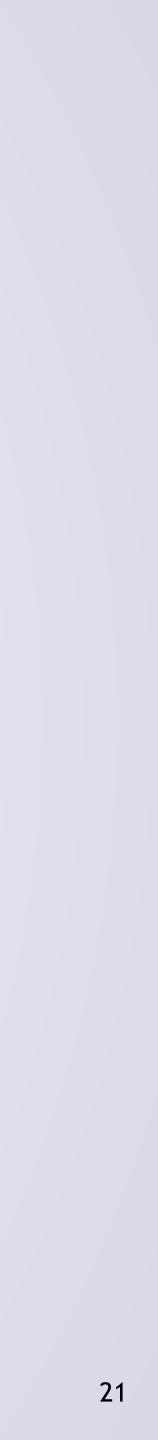
} //thing1 goes out of scope: delete its pointer

Owns a pointer •

When destroyed, deletes owned pointer



std::unique_ptr



std::unique ptr<Thing> thing1 (new Thing); //other code here Thing * tp = thing1.get();

- Owns a pointer
 - When destroyed, deletes owned pointer •
- Can use .get() to get raw pointer



std::unique_ptr

std::unique ptr<Thing> thing1 (new Thing); //other code here Thing * tp = thing1.get(); thing1->doSomething();

- Owns a pointer
 - When destroyed, deletes owned pointer
- Can use .get() to get raw pointer

Can also use * and -> operators



std::unique_ptr

std::unique_ptr

std::unique ptr<Thing> thing1 (new Thing); //... std::unique ptr<Thing> thing2 (thing1);

Assignment operator/copy constructor transfer ownership



- //thing2 owns pointer, thing1 is empty (holds nullptr)



Thing * foo(int x, char c) { Widget * w = new Widget(x); Gadget * g = new Gadget(c); Thing * t = new Thing(w,g); return t;



Exception Safety

- Which guarantee does this make? A: No Throw
- **B:** Strong
 - C: Basic
 - D: No guarantee

Why? If Thing throws exception, then w and g are leaked!







Thing * foo(int x, char c) { std::unique ptr<Widget> w (new Widget(x)); std::unique ptr<Gadget> g (new Gadget(c)); Thing * t = new Thing(w.get(),g.get()); return t;

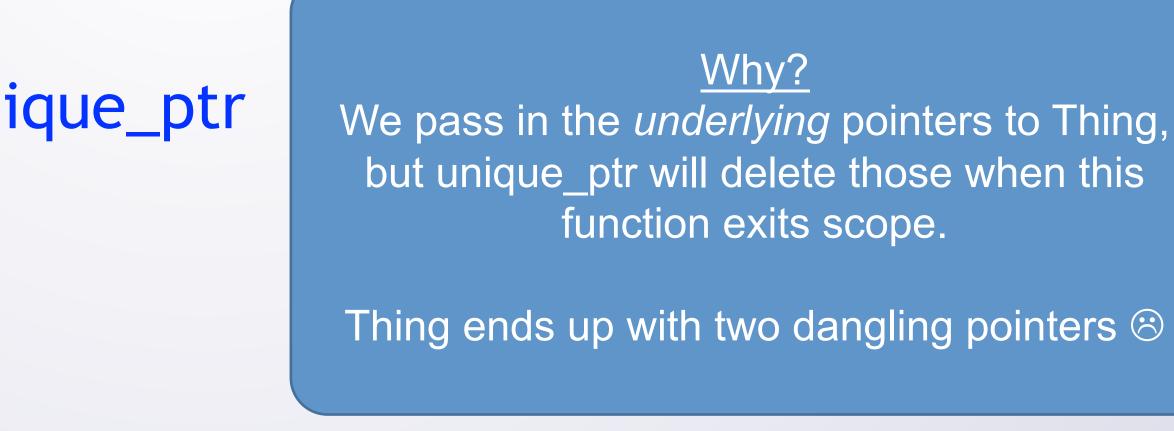
} w and g go out of scope here, so... what happens to their pointers?

Is this code correct?

- A: Yes
- B: No <
- C: I'm lost on unique_ptr



Exception Safety





Thing * foo(int x, char c) { std::unique ptr<Widget> w (new Widget(x)); std::unique ptr<Gadget> g (new Gadget(c)); Thing * t = new Thing(w.release(),g.release()); return t;

release returns the pointer (like get),



Exception Safety

What about this code? No

but also gives up ownership (sets the owned pointer to nullptr)

Why?

What if Thing() constructor throws exception? We still had to run release() first to call it (probably*). Result: more leaked pointers 🛞

* Actually unspecified in the C++ standard.



Thing * foo(int x, char c) { std::unique ptr<Widget> w (new Widget(x)); std::unique ptr<Gadget> g (new Gadget(c)); Thing * t = new Thing(w.release(),g.release()); return t;

What if **new** fails?

"Whether the allocation function is called before evaluating the constructor arguments or after evaluating the constructor arguments but before entering the constructor is unspecified. It is also unspecified whether the arguments to a constructor are evaluated if the allocation function returns the null pointer or exits using an exception."

Duke – C++ standard, 5.3.4 (21)

Exception Safety

What about this code?



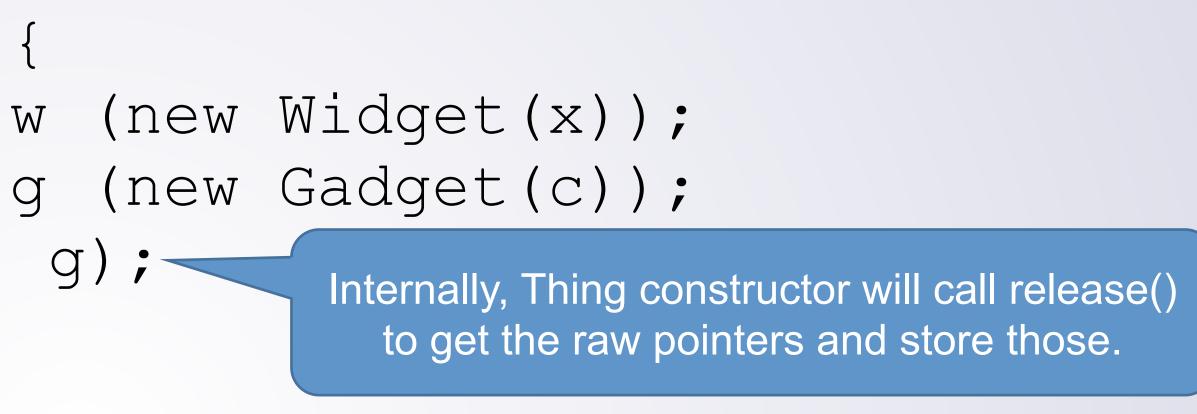
Thing * foo(int x, char c) { std::unique ptr<Widget> w (new Widget(x)); std::unique ptr<Gadget> g (new Gadget(c)); Thing * t = new Thing(w, g); return t;

What am I assuming Thing's constructor takes now? A: Thing (std::unique_ptr<Widget> &, std::unique_ptr<Gadget> &)

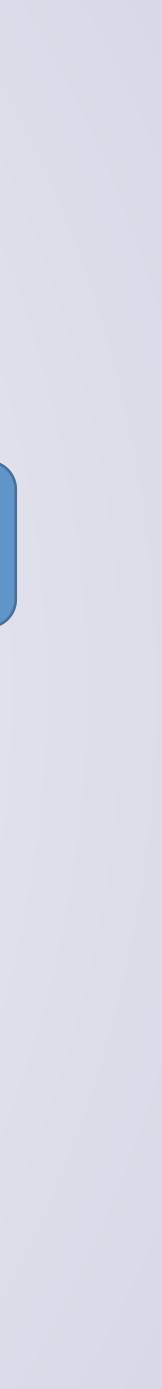
- B: Thing(Widget * , Gadget *)
- C: Thing(Widget, Gadget)
- D: Thing (const Widget & , const Gadget &)



Exception Safety



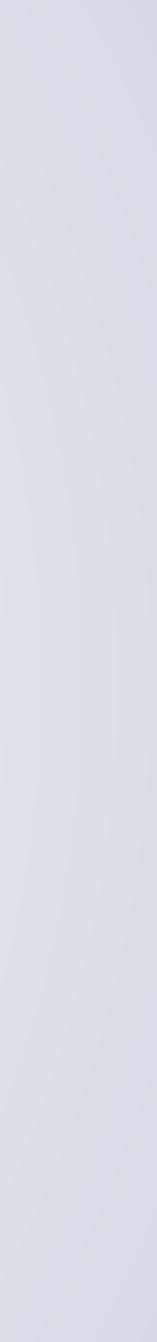
Second: Is this code now correct? A: Yes B: No



Shared Pointers + Weak Pointers

- Unique Pointers: exactly one owner
 - Assignment **transfers** ownership
- Shared Pointers: many owners
 - Copying increments count of owners
 - Destruction decrements counts of owners
 - Object freed when owner count reaches 0
- Weak Pointers: non-owners of shared pointer
 - Can reference object, but does not figure into owner count
 - Use .lock() to obtain shared_ptr: has object (if exists) or nullptr (if not)

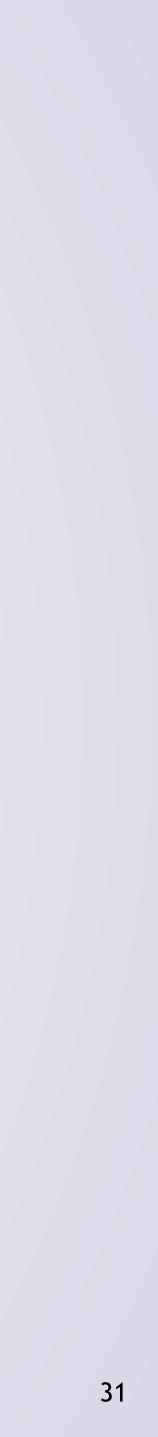




Real C++: Use RAII

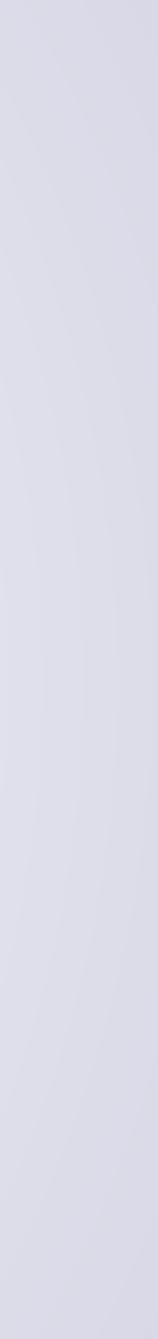
- You learned C++ from C
 - We did a lot of things to transition gently
 - Looked somewhat C-like
 - Less C-like and more C++-like as we progressed
- Real C++:
 - Use RAII for everything





- RAII: C++, but not Java (why not?) •
 - No objects in stack in Java (all in heap...) •
- Java's plan: finally •
 - **ALWAYS** executed, no matter whether exception or not •





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```
public void doAThing(String name) {
  SomeResource sr = null;
  try {
     sr = new SomeResource(name);
     doStuff(sr);
  catch(WhateverException we) {
     dealWithProblem(we);
  finally {
     if(sr != null) {
         sr.close();
```



```
SomeResource sr = null;
try {
   sr = new SomeResource(name);
   doStuff(sr);
finally {
   if(sr != null) {
       sr.close();
```



public void doAThing(String name) throws WhateverException{

Can have try-finally (no catch)

- Allows exception to propagate out
- Cleans up resources

public void doAThing(String name) throws WhateverException{ try (SomeResource sr = new SomeResource(name)) { doStuff(sr);

Java also has try-with-resource



* declare/initialize AutoCloseable object in () after try - can have multiple declarations, separate with; * automatically makes a finally which closes it - closes in reverse order of creation * can have explicit catch or finally if you want

```
SomeResource sr = null;
try {
   sr = new SomeResource(name);
   doStuff(sr);
finally {
   if(sr != null) {
       sr.close();
```

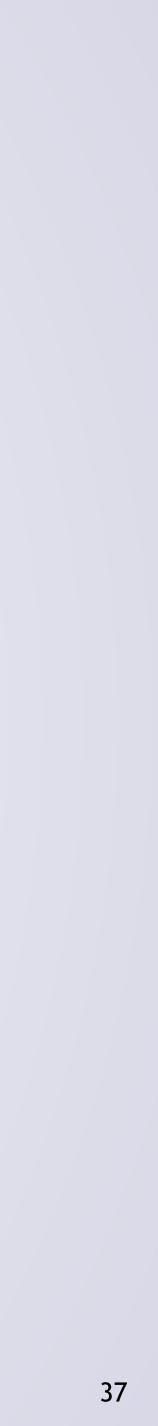


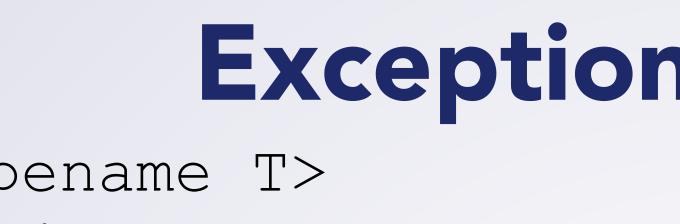
public void doAThing(String name) throws WhateverException{

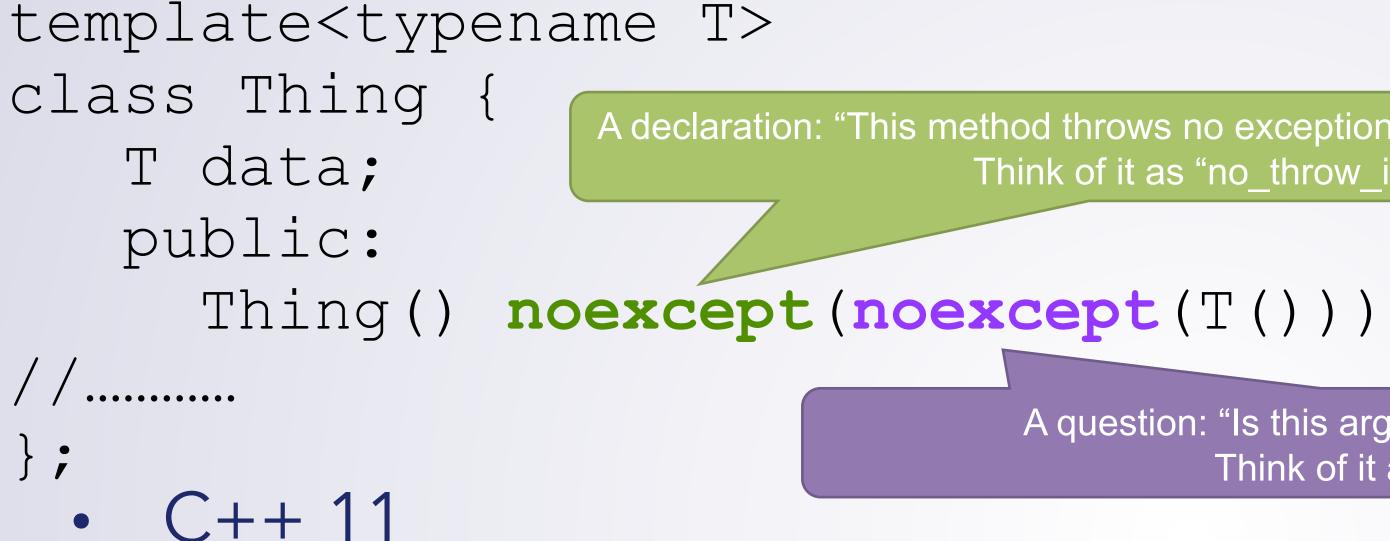
Java's exception specification rules different from C++'s

- C++ 03
 - No declaration: can throw anything
 - Declaration: restricted to those types throw(x, y, z) or throw()
 - Checked at runtime: when exception is thrown
 - If lied, std::unexpected()









- C++03 specifications valid but deprecated •
- noexcept for "no throw"

 - noexcept(expr) queries if expr is declared noexcept •
- If noexcept actually throws, calls std::terminate()

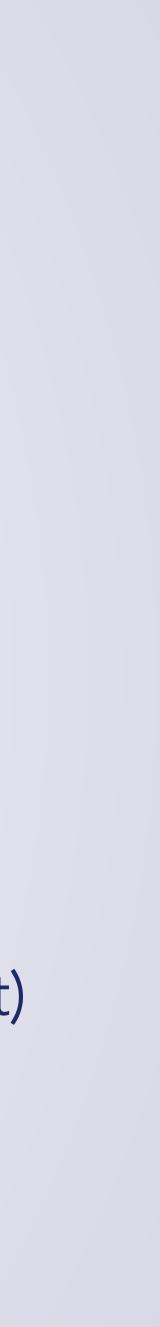


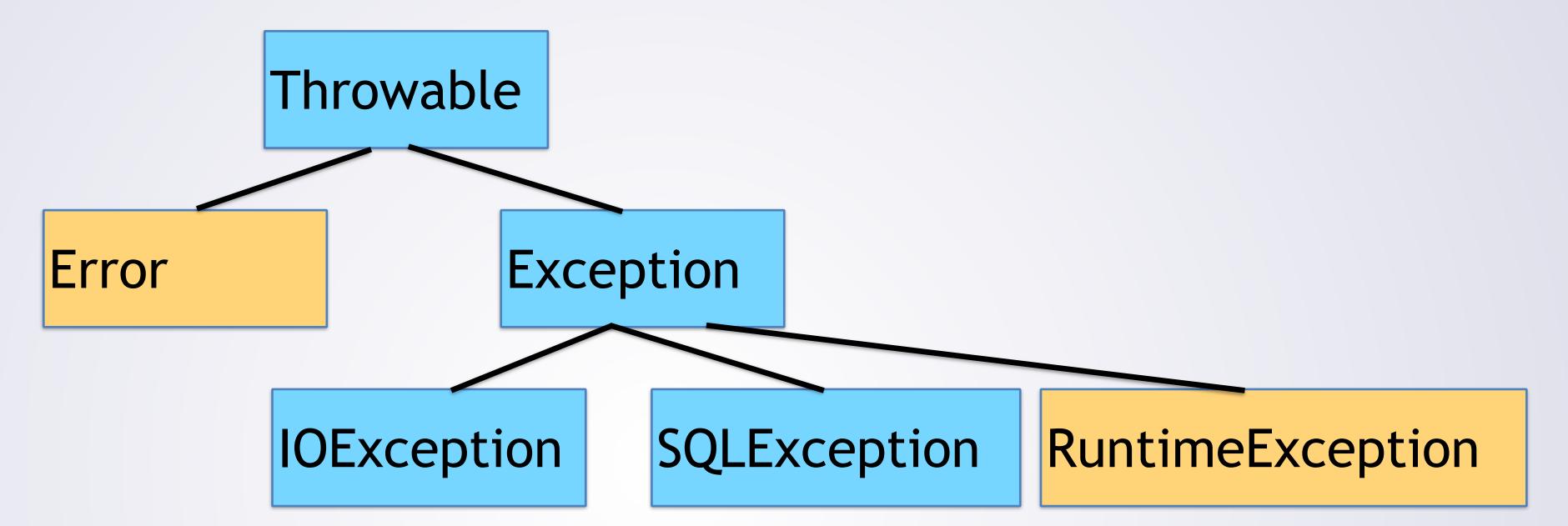
A declaration: "This method throws no exception *if* the argument is true." Think of it as "no_throw_if"

$\{ \}$

A question: "Is this argument marked noexcept?" Think of it as "no_throw?"

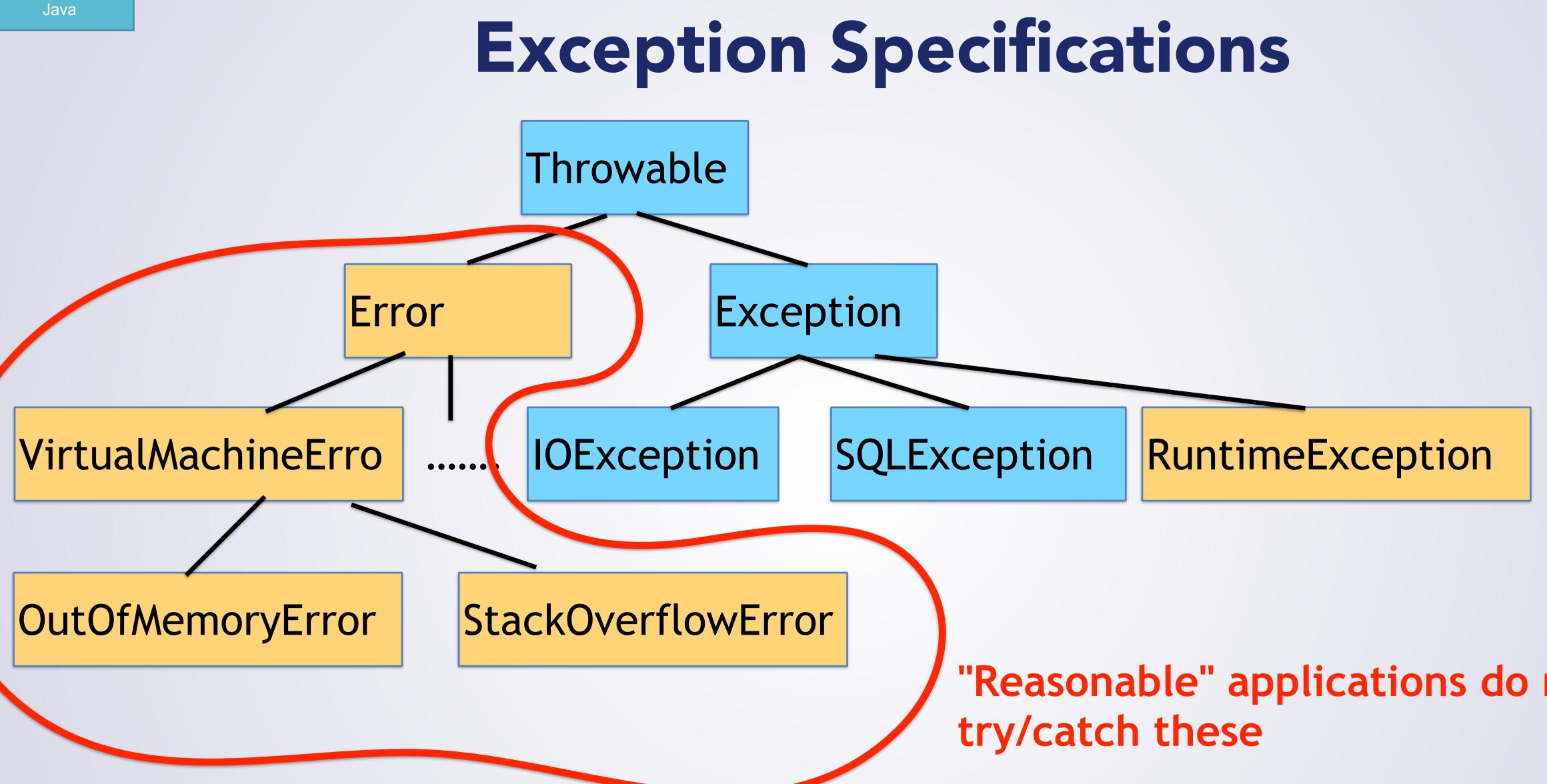
Can take a **boolean expression** to indicate behavior (true=noexcept)





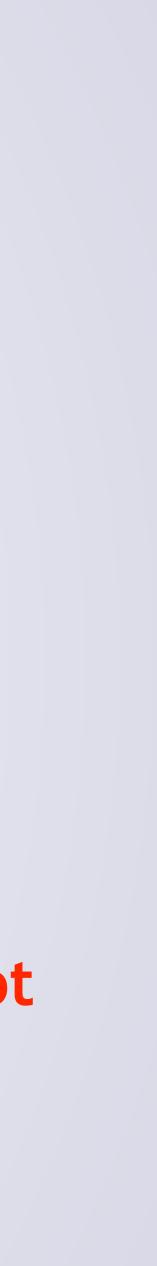
- Java
 - Two types of exceptions: checked and unchecked
 - Checked: exception specifications checked at compile time
 - Compiler ensures you don't lie (aka miss one)
 - Unchecked: no need to declare in spec
 - Possible in too many places, would clutter code



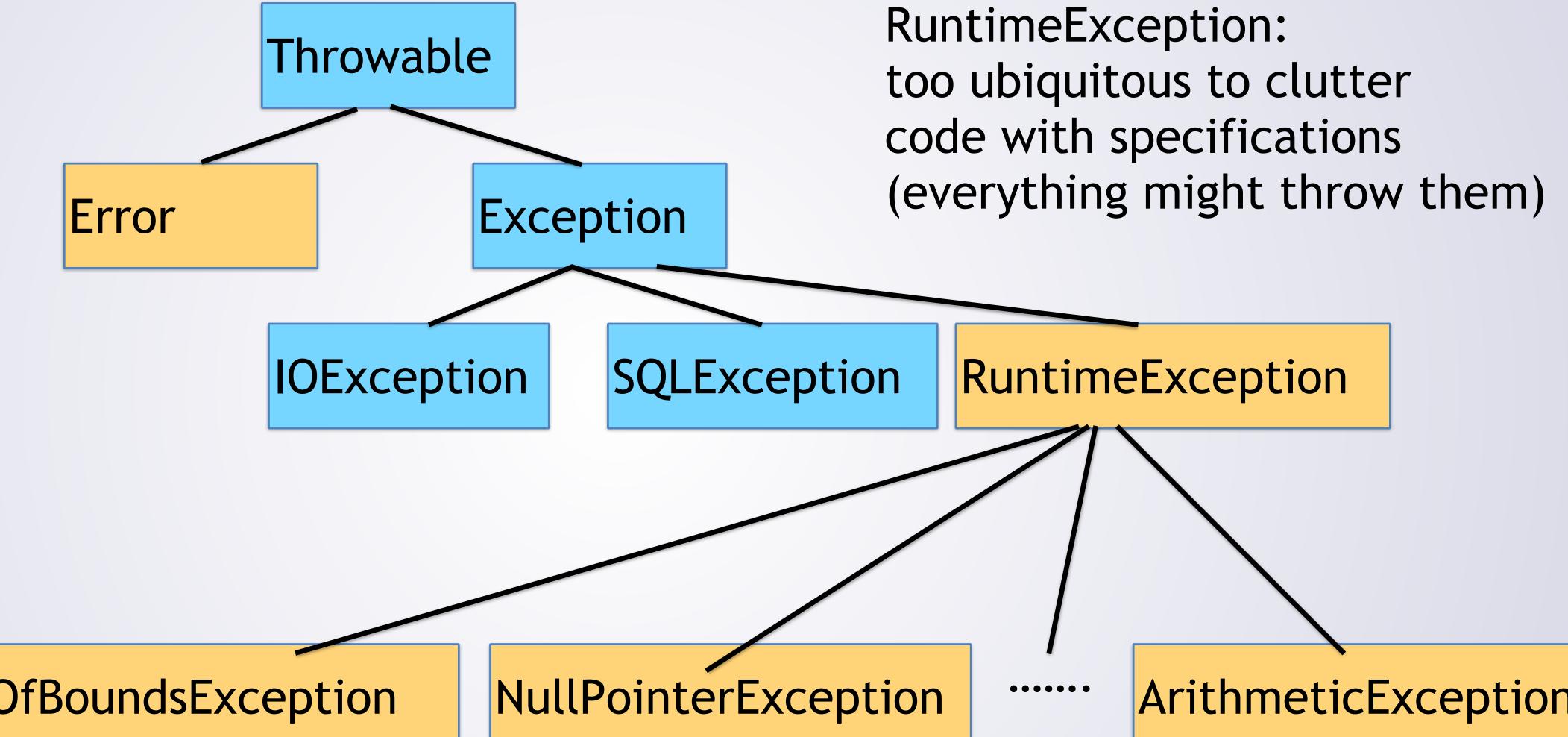




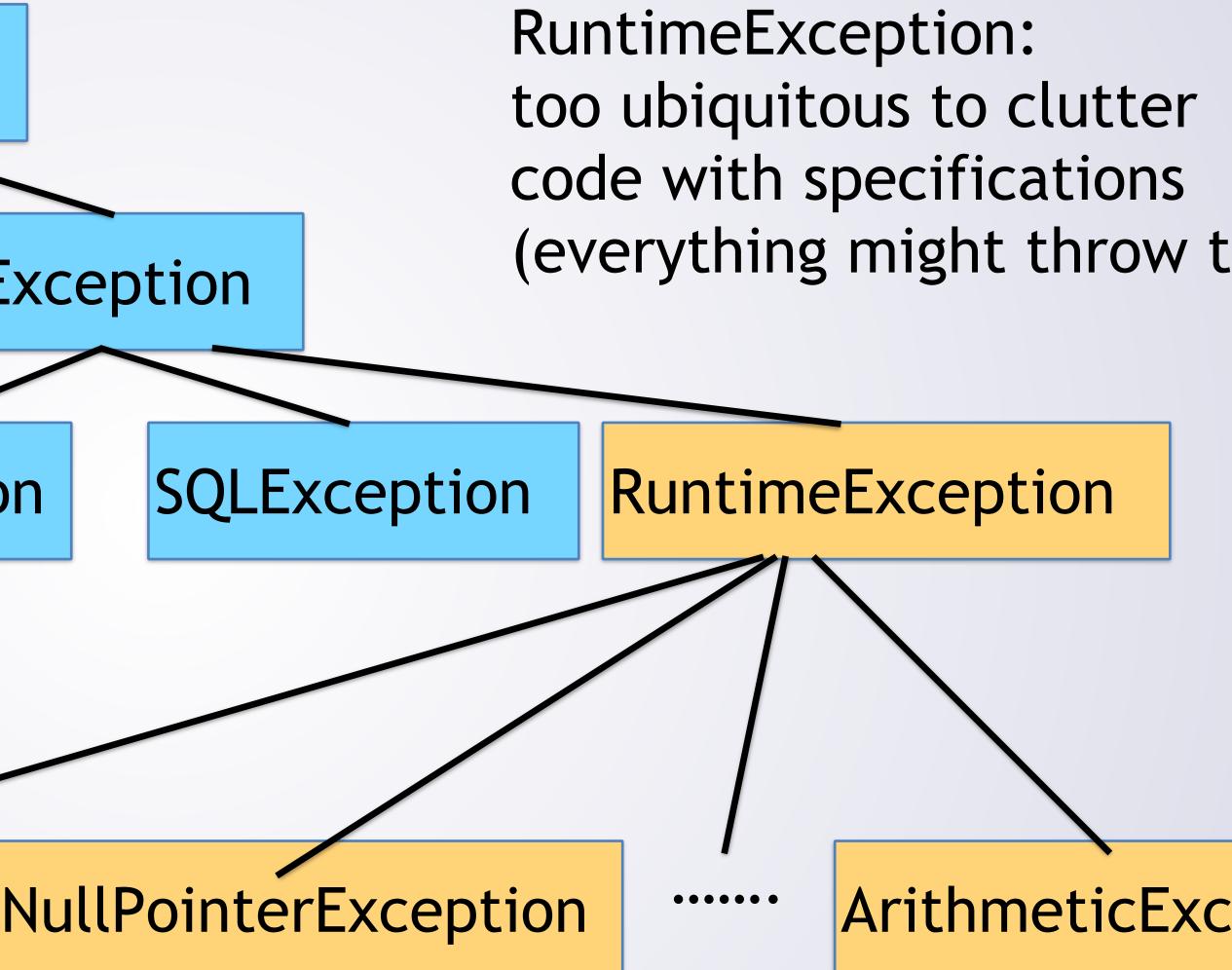
"Reasonable" applications do not



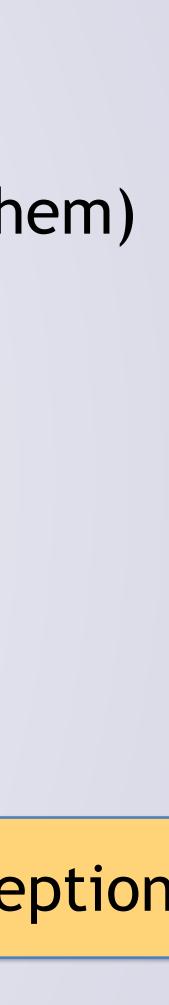




ArrayIndexOutOfBoundsException

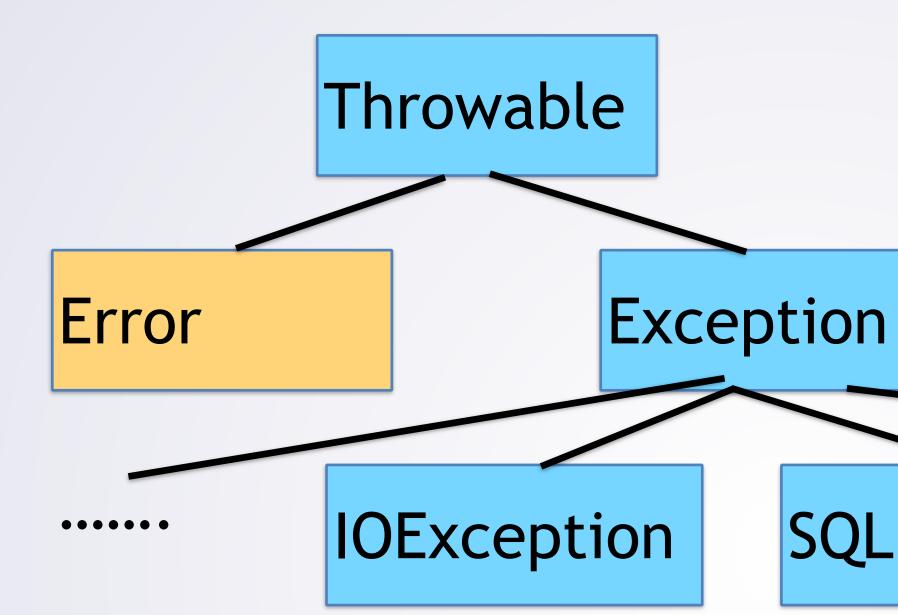






Java

Exception Specifications





Checked exceptions:

- Rare enough to merit specification
- Reasonable enough to try/catch

SQLException

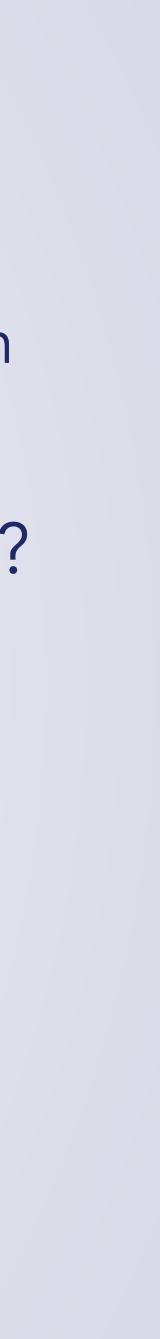
RuntimeException



Java: Finalizers

- Java objects have .finalize()
 - "Called by the garbage collector on an object when garbage collection determines that there are no more references to the object."
- Seems like maybe we could use this to help resource management?

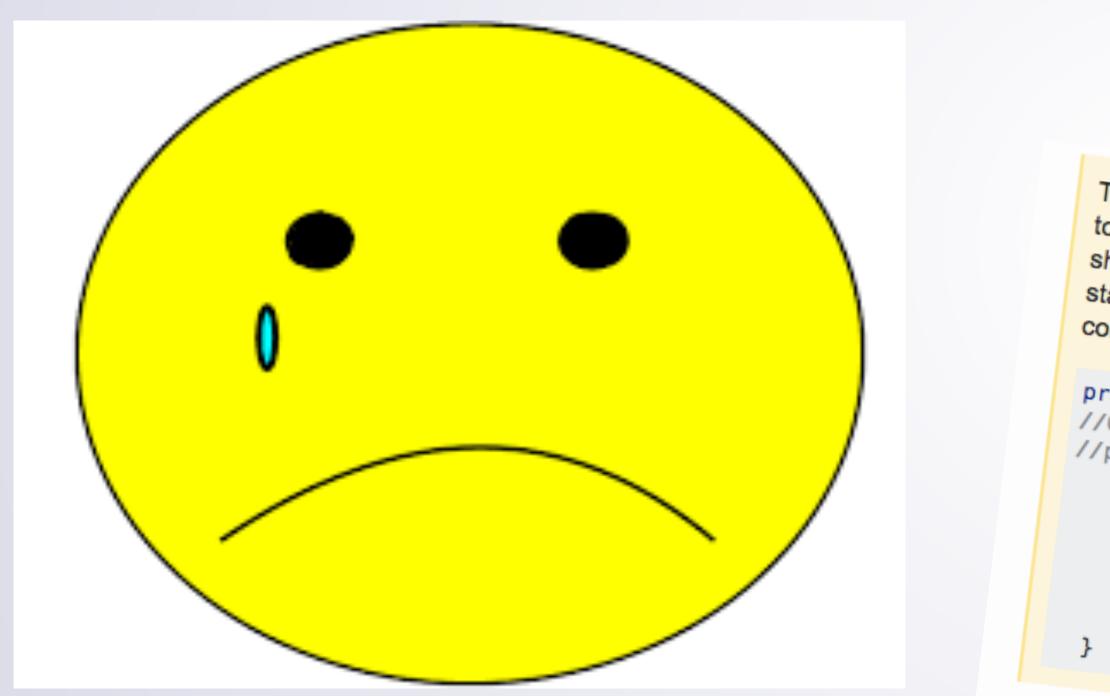




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Lets Look at Stack Overflow

When the IO resource is an instance variable, then you should close it in the finalize() method.





http://stackoverflow.com/questions/12958440/closing-class-io-resources-in-overridden-finalize-method http://stackoverflow.com/questions/8051863/how-can-i-close-the-socket-in-a-proper-way

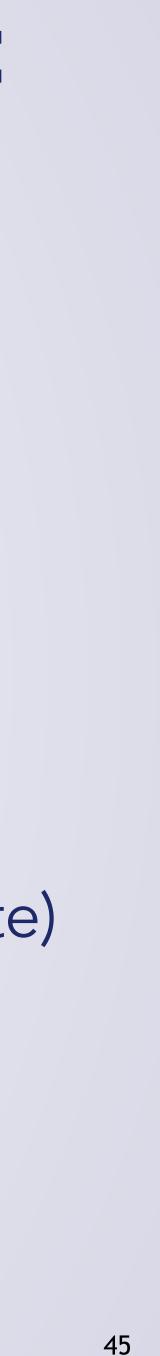
The finalize() method is called by the Java virtual machine (JVM) before the program exits to give the program a chance to clean up and release resources. Multi-threaded programs should close all Files and Sockets they use before exiting so they do not face resource starvation. The call to server.close() in the finalize() method closes the Socket protected void finalize(){ //Objects created in run method are finalized when //program terminates and thread exits server.close(); } catch (IOException e) { System.out.println("Could not close socket");



Finalizer: NOT For Resource Management

- Do NOT try to use finalizers for resource management!
 - No guarantee of when they will run (may never gc object!) •
- Do NOT use finalizers in general
 - May run on other threads (possibly multiple finalizers at once) •
 - Were you thinking about how to synchronize them?
 - What about deadlock?
 - Likely to run when memory is scarce (may cause problems if you allocate) Could accidentally make object re-referenceable? •







- Handling problems: exceptions
- C++
 - temp-and-swap
 - RAII
 - Smart Pointers
- Java
 - finally •
 - specifications
 - finalizers (and why they are not what you need for this)



Exceptions