

The Planted Matching Problem

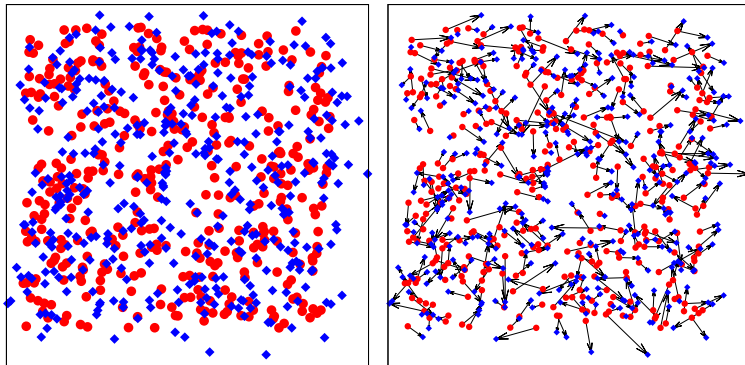
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Joint work with
Mehrdad Moharrami (UIUC) and Cristopher Moore (Santa Fe Institute)

October 17, 2022
INFORMS Annual Meeting

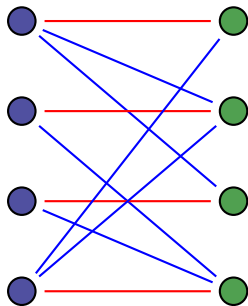
Motivating application: particle tracking



[Chertkov-Kroc-Krzakala-Vergassola-Zdeborová PNAS'10]

- Tracking particles advected by turbulent fluid flow
- **Goal:** recover the latent correspondence between particles

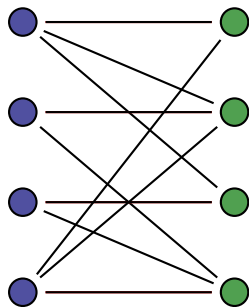
The planted matching model



- A complete bipartite graph
- A hidden perfect matching M^*
- Edge weight

$$W_e \stackrel{\text{ind.}}{\sim} \begin{cases} P & e \in M^* \\ Q & e \notin M^* \end{cases}$$

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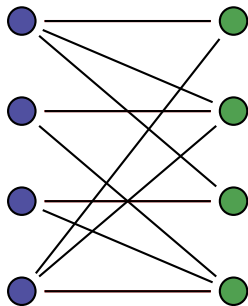


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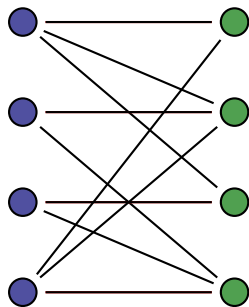


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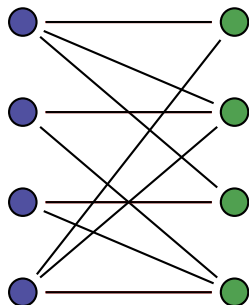


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- Minimum-weight matching is Maximum Likelihood Estimator \hat{M}_{ML}
- How much does \hat{M}_{ML} have in common with M^* ?

Main result: phase transition at $\lambda = 4$

Theorem (Moharrami-Moore-X. AAP'21)

$$\text{overlap: } \lim_{n \rightarrow \infty} \frac{1}{n} \mathbb{E} \left[\left| \widehat{M}_{\text{ML}} \cap M^* \right| \right] = \begin{cases} 1 & \text{if } \lambda \geq 4 \\ \alpha(\lambda) & \text{if } 0 < \lambda < 4 \end{cases}$$

where $\alpha(\lambda) = 1 - 2 \int_0^\infty (1 - F(x))(1 - G(x)) V(x)W(x) dx$,

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and F, G, V, W is the unique solution to the ODE system

$$\dot{F} = (1 - F)(1 - G)V$$

$$\dot{G} = -(1 - F)(1 - G)W$$

$$\dot{V} = \lambda(V - F)$$

$$\dot{W} = -\lambda(W - G)$$

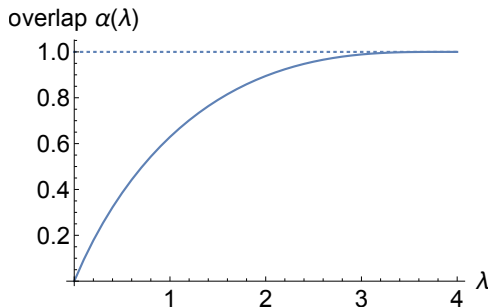
Boundary conditions: $F(x), V(x), G(-x), W(-x) \rightarrow \begin{cases} 1 & x \rightarrow +\infty \\ 0 & x \rightarrow -\infty \end{cases}$

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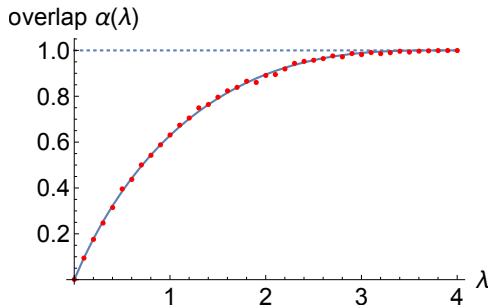


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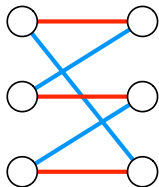
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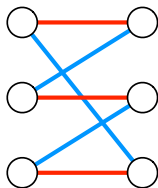
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When $\lambda \geq 4$: count augmenting cycles



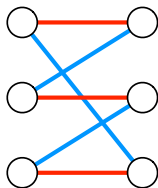
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- Probability that an alternating cycle of length 2ℓ is **augmenting**:

$$\mathbb{P}[\text{Erlang}(\ell, \lambda) \geq \text{Erlang}(\ell, 1/n)] \leq \left(\frac{\lambda n}{4}\right)^{-\ell}$$

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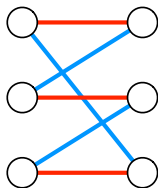


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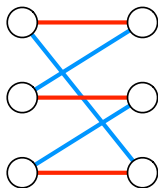


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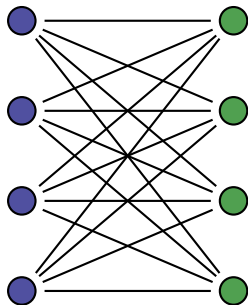


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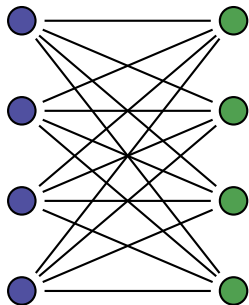
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- \Rightarrow Expected # of such augmenting cycles is at most $(\lambda/4)^{-\ell} e^{-\ell^2/2n}$
- \Rightarrow Sum over $\ell \implies \mathbb{E} \left[\left| \widehat{M}_{\text{ML}} \Delta M^* \right| \right] = o(n)$

Warmup: the (un-planted) random assignment problem



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- Weights are i.i.d. $\text{Exp}(1/n)$
- Cost of minimum matching?

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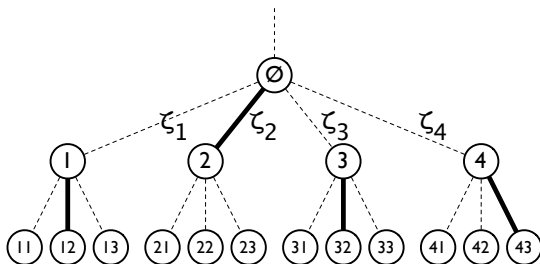
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[Walkup'79, Mézard-Parisi'87, Steele'97, Aldous'01, Nair-Prabhakar-Sharma'05, Wästlund'09]

$$\mathbb{E} \left[\min_{M \in \mathcal{M}} \frac{1}{n} \sum_{e \in M} W_e \right] = 1 + \frac{1}{4} + \frac{1}{9} + \cdots + \frac{1}{n^2} \rightarrow \frac{\pi^2}{6}$$

Poisson-weighted infinite tree approximation

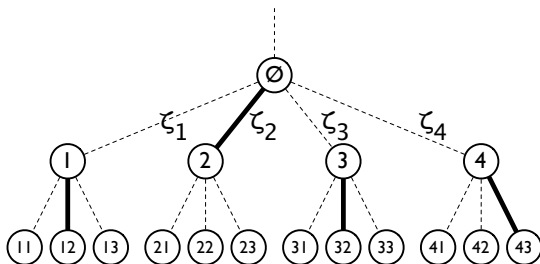
Cavity method: model as a tree [Mézard-Parisi '87, Aldous'00]



sort edge weights $W_{\emptyset,1}, W_{\emptyset,2}, \dots$ from smallest to largest:
arrivals ζ_1, ζ_2, \dots of a Poisson process with rate 1

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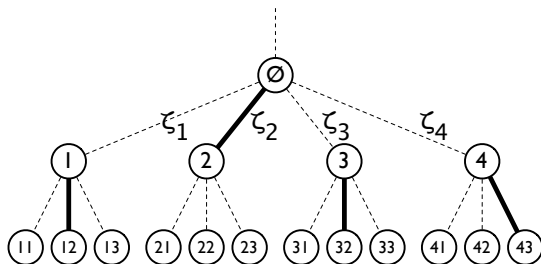


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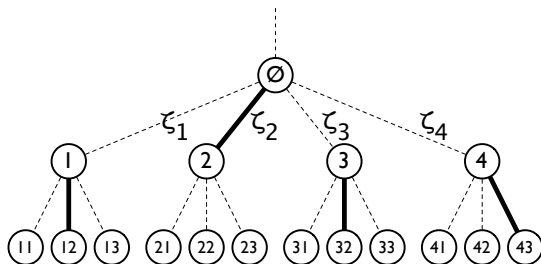
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$$X_{\emptyset} = \min_{i \geq 1} \{W_{\emptyset,i} - X_i\} \implies X \stackrel{d}{=} \min_{i \geq 1} \{\zeta_i - X_i\}$$

From distributional to differential equations

$X \stackrel{d}{=} \min \{\zeta_i - X_i\}$ where ζ_i are Poisson arrivals

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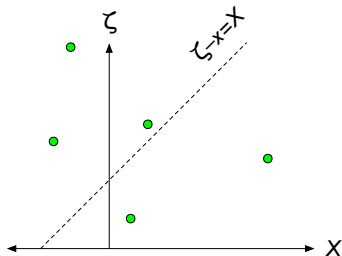
Define the ccdf $\bar{F}(x) = 1 - F(x) = \mathbb{P}[X > x] = \mathbb{P}[\forall i : \zeta_i - x > X_i]$

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Generate pairs (ζ_i, X_i) : two-dimensional Poisson process with density F'

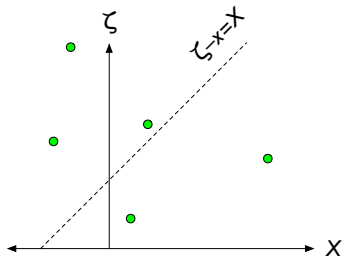


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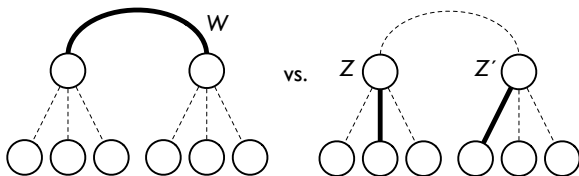
$$\bar{F}(x) = \exp\left(-\int_{-x}^{\infty} \bar{F}(t) dt\right) \Rightarrow \frac{dF(x)}{dx} = F(x)F(-x)$$

From distributional to differential equations, cont'd

$$\frac{dF(x)}{dx} = F(x)F(-x) \implies F(x) = \frac{e^x}{1 + e^x}$$

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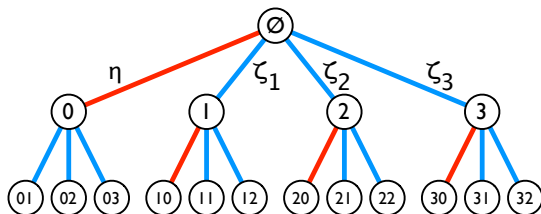


Contribution of a single edge:

$$\int_0^\infty w \mathbb{P}[Z + Z' \geq w] dw = \frac{1}{4} \text{Var}[Z + Z'] = \frac{1}{2} \text{Var}[Z] = \frac{\pi^2}{6}$$

Planted Poisson-weighted infinite tree

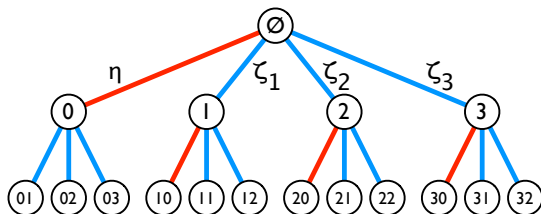
Partner in planted matching is either parent or child 0, other children sorted 1, 2, 3, ...



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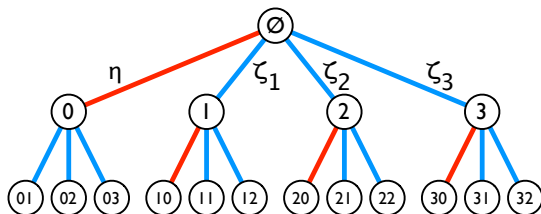
Recursion:

$$X_{\emptyset} = \min \left\{ W_{\emptyset,0} - X_0, \min_{i \geq 1} \{W_{\emptyset,i} - X_i\} \right\}$$

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From distributional to differential equations, redux

$$\begin{aligned} Y &\stackrel{d}{=} \min \{ \eta - Z, Z' \} \\ Z &\stackrel{d}{=} \min \{ \zeta_i - Y_i \}_{i=1}^{\infty} \end{aligned}$$

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\dot{V} and \dot{W} from $\eta \sim \text{Exp}(\lambda)$, integration by parts

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Phase transition of ODE at $\lambda = 4$

$$\begin{aligned}\dot{F} &= (1 - F)(1 - G)V \\ \dot{G} &= -(1 - F)(1 - G)W \\ \dot{V} &= \lambda(V - F) \\ \dot{W} &= -\lambda(W - G)\end{aligned}$$

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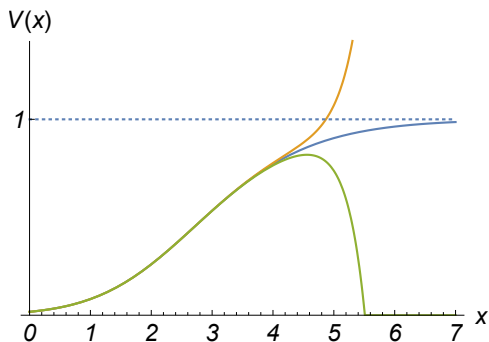
Lemma

There is a unique solution if and only if $\lambda < 4$.

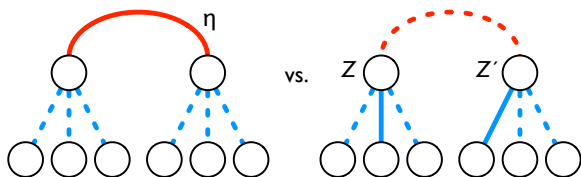
Geometric interpretation of uniqueness

When $\lambda < 4$, $(F = 1, G = 0, V = 1, W = 0)$ is a **saddle point**:

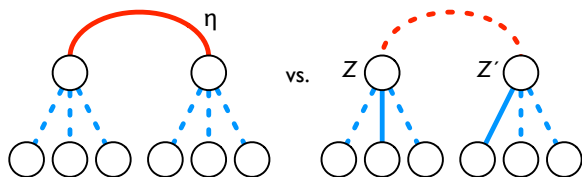
There exists a unique initial condition from which we approach the saddle along its unstable manifold



Finally, computing the overlap for $\lambda < 4$

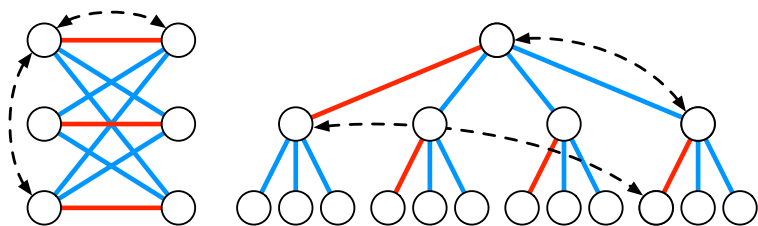


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$$\begin{aligned}
 \alpha(\lambda) &= \mathbb{P}[\eta < Z + Z'] = 1 - \mathbb{E}_\eta \int_{-\infty}^{+\infty} f(x)F(\eta - x) dx \\
 &= 1 - \int_{-\infty}^{+\infty} f(x) \mathbb{E}_\eta F(\eta - x) dx \\
 &= 1 - \int_{-\infty}^{+\infty} (1 - F(x))(1 - G(x))V(x)W(x) dx \\
 &= 1 - 2 \int_0^{+\infty} (1 - F(x))(1 - G(x))V(x)W(x) dx
 \end{aligned}$$

Proving it: Local weak convergence (Aldous 1992, 2001)



- Construct a *spatially invariant* M_{opt} on T_{∞} using message passing
- Show $(K_{n,n}, M_{\text{min}})$ converges locally to $(T_{\infty}, M_{\text{opt}})$
 - ▶ Local treelikeness of light edges
 - ▶ Almost-doubly-stochastic matrix

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- Sharp threshold for almost perfect recovery: $\lambda = 4$
- Exact characterization of overlap of MLE by system of ODEs

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Extensions

- $\lambda = 4$ is the sharp information-theoretic threshold [Ding-Wu-X.-Yang '21]:

$\lambda = 4 - \epsilon$: Optimal reconstruction error is $\exp(-\Theta(1/\sqrt{\epsilon}))$

- Sharp threshold in general (d, P, Q) model [Ding-Wu-X.-Yang '21]:

$$\sqrt{d} \int \sqrt{dP dQ} = 1$$

- Non-IID spatial model [Kunisky-Niles-Weed '22, Wang-Wu-X.-Yolou '22]

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Reference

- M. Moharrami, C. Moore, & J. Xu, *The planted matching problem: Phase transitions and exact results*. *Annals of Applied Probability*, 2021.
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